
Examples from “A Decision Procedure for Probability Calculus with Applications”

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First, load the **PrSAT** package (freely downloadable from the **PrSAT** website: <http://fitelson.org/PrSAT/>).

<< **PrSAT**^

■ Example #1: A probability model in which three statements **A**, **B**, and **C** are *pariwise* independent but not *mutually* independent

Note: the additional equational side-constraints [**Pr**[**A**] == $\frac{1}{10}$, **Pr**[**B**] == $\frac{1}{10}$, **Pr**[**C**] == $\frac{1}{10}$] are added here to reduce the number of variables of the problem, so that the model is found *much* more quickly. This is a very useful heuristic for speeding-up model searches. The option **Probabilities→Regular** guarantees that the model generated is *regular* (i.e., that it assigns nonzero probability to all state descriptions of the minimal sentential language required for the expression of the problem given). Finally, the option **BypassSearch→True** tells *Mathematica* to skip Blum’s random search add-on, and send the problem straight to the decision procedure.

```

MODEL1 = PrSAT [
  {
    Pr[A ∧ B ∧ C] ≠ Pr[A] Pr[B] Pr[C] ,
    Pr[A ∧ B] == Pr[A] Pr[B] ,
    Pr[A ∧ C] == Pr[A] Pr[C] ,
    Pr[B ∧ C] == Pr[B] Pr[C] ,

    (* Heuristic -- add additional
      equational side-constraints *)
    Pr[A] ==  $\frac{1}{10}$  , Pr[B] ==  $\frac{1}{10}$  , Pr[C] ==  $\frac{1}{10}$ 
  } ,
  Probabilities → Regular ,
  BypassSearch → True
]

{ {A → {a2 , a5 , a6 , a8} ,
  B → {a3 , a5 , a7 , a8} , C → {a4 , a6 , a7 , a8} ,
  Ω → {a1 , a2 , a3 , a4 , a5 , a6 , a7 , a8} } ,
  { a1 →  $\frac{1459}{2000}$  , a2 →  $\frac{161}{2000}$  , a3 →  $\frac{161}{2000}$  , a4 →  $\frac{161}{2000}$  ,
    a5 →  $\frac{19}{2000}$  , a6 →  $\frac{19}{2000}$  , a7 →  $\frac{19}{2000}$  , a8 →  $\frac{1}{2000}$  } }

```

PrSAT also includes a **TruthTable** function, which allows for the visualization of a model, as a stochastic truth-table:

TruthTable [MODEL1]

A	B	C	var	Pr
T	T	T	a ₈	$\frac{1}{2000}$
T	T	F	a ₅	$\frac{19}{2000}$
T	F	T	a ₆	$\frac{19}{2000}$
T	F	F	a ₂	$\frac{161}{2000}$
F	T	T	a ₇	$\frac{19}{2000}$
F	T	F	a ₃	$\frac{161}{2000}$
F	F	T	a ₄	$\frac{161}{2000}$
F	F	F	a ₁	$\frac{1459}{2000}$

If we set **BypassSearch**→**False**, then Blum's random-search add-on is consulted first, and a model is (usually) found (relatively) quickly *even without any side-constraints!*

```

MODEL11 = PrSAT[
  {
    Pr[A ∧ B ∧ C] ≠ Pr[A] Pr[B] Pr[C],
    Pr[A ∧ B] == Pr[A] Pr[B],
    Pr[A ∧ C] == Pr[A] Pr[C],
    Pr[B ∧ C] == Pr[B] Pr[C]
  },
  Probabilities → Regular,
  BypassSearch → False
]

```

$$\left\{ \begin{array}{l}
 \mathbf{A} \rightarrow \{a_2, a_5, a_6, a_8\}, \\
 \mathbf{B} \rightarrow \{a_3, a_5, a_7, a_8\}, \mathbf{C} \rightarrow \{a_4, a_6, a_7, a_8\}, \\
 \mathbf{\Omega} \rightarrow \{a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8\}, \\
 a_1 \rightarrow \frac{84\,418 - 39\sqrt{4\,676\,097}}{56\,277}, \\
 a_2 \rightarrow \frac{-42\,296 + 39\sqrt{4\,676\,097}}{168\,831}, \\
 a_3 \rightarrow \frac{-42\,296 + 39\sqrt{4\,676\,097}}{168\,831}, \\
 a_4 \rightarrow \frac{-42\,296 + 39\sqrt{4\,676\,097}}{168\,831}, a_5 \rightarrow \frac{1}{999}, \\
 a_6 \rightarrow \frac{1}{999}, a_7 \rightarrow \frac{1}{999}, a_8 \rightarrow \frac{42}{169} \}
 \end{array} \right.$$

TruthTable [MODEL11]

A	B	C	var	Pr
T	T	T	a ₈	$\frac{42}{169}$
T	T	F	a ₅	$\frac{1}{999}$
T	F	T	a ₆	$\frac{1}{999}$
T	F	F	a ₂	$\frac{-42\,296+39\sqrt{4\,676\,097}}{168\,831}$
F	T	T	a ₇	$\frac{1}{999}$
F	T	F	a ₃	$\frac{-42\,296+39\sqrt{4\,676\,097}}{168\,831}$
F	F	T	a ₄	$\frac{-42\,296+39\sqrt{4\,676\,097}}{168\,831}$
F	F	F	a ₁	$\frac{84\,418-39\sqrt{4\,676\,097}}{56\,277}$

We can also Evaluate probabilities on given models, as follows:

EvaluateProbability[
{Pr[A | B], Pr[B | A ∧ ¬ C]}, MODEL11]

$$\left\{ \frac{42\,127}{168\,831 \left(\frac{42\,296}{168\,831} + \frac{-42\,296+39\sqrt{4\,676\,097}}{168\,831} \right)}, \frac{1}{999 \left(\frac{1}{999} + \frac{-42\,296+39\sqrt{4\,676\,097}}{168\,831} \right)} \right\}$$

% // N

{0.499521, 0.00400401}

■ A Simultaneous Countermodel to the **S**-instances of both (*) and (†)

Again, the additional equational side-constraints $[\mathbf{Pr}[\mathbf{H}] == \frac{1}{2}, \mathbf{Pr}[\mathbf{E1}] == \frac{1}{4}, \mathbf{Pr}[\mathbf{E2}] == \frac{3}{4}]$ are added to speed the model search, and **Probabilities** → **Regular** indicates that we are asking **PrSAT** to find a *regular* probability model.

```

MODEL2 = PrSAT [
  {
    Pr[H | E1] > Pr[H],
    Pr[H | E2] > Pr[H],
    Pr[H | E1] > Pr[H | E2],
    Pr[H | E1] - Pr[H | ¬ E1] <
      Pr[H | E2] - Pr[H | ¬ E2],
    Pr[H | E1 ∧ E2] - Pr[H | ¬ E1 ∧ E2] ==
      Pr[H | E1] - Pr[H | ¬ E1],
    Pr[H | E2 ∧ E1] - Pr[H | ¬ E2 ∧ E1] ==
      Pr[H | E2] - Pr[H | ¬ E2],
    Pr[H | E1 ∧ E2] - Pr[H | ¬ (E1 ∧ E2)] <
      Pr[H | E2] - Pr[H | ¬ E2],

    (* Heuristic -- add additional
       equational side-constraints *)
    Pr[H] ==  $\frac{1}{2}$ , Pr[E1] ==  $\frac{1}{4}$ , Pr[E2] ==  $\frac{3}{4}$ 
  },
  Probabilities → Regular,
  BypassSearch → True
]

{ {E1 → {a2, a5, a6, a8},
  E2 → {a3, a5, a7, a8}, H → {a4, a6, a7, a8},
  Ω → {a1, a2, a3, a4, a5, a6, a7, a8}},
  { a1 →  $\frac{147}{1024}$ , a2 →  $\frac{193}{5120}$ , a3 →  $\frac{1341}{5120}$ , a4 →  $\frac{45}{1024}$ ,
    a5 →  $\frac{291}{5120}$ , a6 →  $\frac{127}{5120}$ , a7 →  $\frac{1539}{5120}$ , a8 →  $\frac{669}{5120}$  } }

```

TruthTable [MODEL2]

E1	E2	H	var	Pr
T	T	T	a ₈	$\frac{669}{5120}$
T	T	F	a ₅	$\frac{291}{5120}$
T	F	T	a ₆	$\frac{127}{5120}$
T	F	F	a ₂	$\frac{193}{5120}$
F	T	T	a ₇	$\frac{1539}{5120}$
F	T	F	a ₃	$\frac{1341}{5120}$
F	F	T	a ₄	$\frac{45}{1024}$
F	F	F	a ₁	$\frac{147}{1024}$

■ A Simultaneous Countermodel to two claims concerning Hawthorne & Fitelson's new Bayesian approach to the raven paradox

The following single model shows that neither of the following two claims:

$$(6) \Pr(H \mid \sim R \ \& \ \sim B) > \Pr(H)$$

$$(7) \Pr(H \mid \sim R \ \& \ B) < \Pr(H)$$

follows from the following three claims:

$$(1) \Pr(R \mid H \ \& \ B) = 1$$

$$(2) \Pr(\sim B) > \Pr(R)$$

$$(C) \Pr(H \mid R) \geq \Pr(H \mid \sim B)$$

Here, a regular model is impossible (since one of the constraints requires a zero probability for one of the state descriptions). But, by adding the constraint $\Pr [(\neg H) \wedge B \wedge (\neg R)] > 0$, we can ensure that this is the only zero in the model. And, as usual, we add equational side-constraints $[\Pr [H] = \frac{60}{100}, \Pr [R] = \frac{20}{100}, \Pr [B] = \frac{10}{100}]$ to speed the model-finding process by reducing the number of free variables in the problem.

MODEL3 = PrSAT [

{

Pr[**H** \wedge **R** \wedge (\neg **B**)] == 0,

Pr[(\neg **H**) \wedge **B** \wedge (\neg **R**)] > 0,

Pr[\neg **B**] > **Pr**[**R**],

Pr[**H** | **R**] == **Pr**[**H** | \neg **B**],

Pr[**H** | (\neg **R**) \wedge (\neg **B**)] < **Pr**[**H**],

Pr[**H** | **B** \wedge (\neg **R**)] > **Pr**[**H**]

},

Constraints \rightarrow

{**Pr**[**B**] == $\frac{1}{10}$, **Pr**[**H**] == $\frac{1}{2}$, **Pr**[**R**] == $\frac{1}{10}$ },

BypassSearch \rightarrow **True**

]

{**B** \rightarrow {**a**₂, **a**₅, **a**₆, **a**₈},

H \rightarrow {**a**₃, **a**₅, **a**₇, **a**₈}, **R** \rightarrow {**a**₄, **a**₆, **a**₇, **a**₈},

Ω \rightarrow {**a**₁, **a**₂, **a**₃, **a**₄, **a**₅, **a**₆, **a**₇, **a**₈}},

{**a**₁ \rightarrow $\frac{141}{320}$, **a**₂ \rightarrow $\frac{21}{2560}$, **a**₃ \rightarrow $\frac{225}{512}$, **a**₄ \rightarrow $\frac{51}{2560}$,

a₅ \rightarrow $\frac{3}{256}$, **a**₆ \rightarrow $\frac{1}{32}$, **a**₇ \rightarrow 0, **a**₈ \rightarrow $\frac{25}{512}$ }}

TruthTable [**MODEL3**]

B	H	R	var	Pr
T	T	T	a ₈	$\frac{25}{512}$
T	T	F	a ₅	$\frac{3}{256}$
T	F	T	a ₆	$\frac{1}{32}$
T	F	F	a ₂	$\frac{21}{2560}$
F	T	T	a ₇	0
F	T	F	a ₃	$\frac{225}{512}$
F	F	T	a ₄	$\frac{51}{2560}$
F	F	F	a ₁	$\frac{141}{320}$

We can also solve this one (quickly) with Blum's random search add-on, and with no side-constraints:

```

PrSAT [
  {
    Pr[H  $\wedge$  R  $\wedge$  ( $\neg$  B)] == 0,
    Pr[ ( $\neg$  H)  $\wedge$  B  $\wedge$  ( $\neg$  R)] > 0,
    Pr[ $\neg$  B] > Pr[R],
    Pr[H | R] == Pr[H |  $\neg$  B],
    Pr[H | ( $\neg$  R)  $\wedge$  ( $\neg$  B)] < Pr[H],
    Pr[H | B  $\wedge$  ( $\neg$  R)] > Pr[H]
  },
  BypassSearch  $\rightarrow$  False
]

{ {B  $\rightarrow$  {a2, a5, a6, a8},
  H  $\rightarrow$  {a3, a5, a7, a8}, R  $\rightarrow$  {a4, a6, a7, a8},
   $\Omega$   $\rightarrow$  {a1, a2, a3, a4, a5, a6, a7, a8}},
  {a1  $\rightarrow$   $\frac{36\,693}{94\,240}$ , a2  $\rightarrow$   $\frac{2}{15}$ , a3  $\rightarrow$  0, a4  $\rightarrow$   $\frac{1}{93}$ ,
  a5  $\rightarrow$   $\frac{11}{32}$ , a6  $\rightarrow$   $\frac{7}{57}$ , a7  $\rightarrow$  0, a8  $\rightarrow$  0}}

```

■ Theorems and Countermodels from Sobel's "Lotteries and Miracles"

Since Sobel's problems only involve two atomic sentences, no heuristics are needed to yield a fast solution by the decision procedure (and Blum's random search add-on is also not necessary, since the decision procedure is quite fast in such cases).

■ A PrSAT model showing that Sobel's (1)-(3) do *not* entail (4)

$$\text{MODEL4} = \text{PrSAT} \left[\begin{array}{l} \left\{ \text{Pr}[\mathbf{T}] < \frac{1}{2}, \right. \\ \text{Pr}[\mathbf{T} \mid \mathbf{W}] > \frac{1}{2}, \\ \text{Pr}[\mathbf{W} \mid \mathbf{T}] > \frac{1}{2}, \\ \left. \text{Pr}[\mathbf{T} \mid \neg \mathbf{W}] > \text{Pr}[\mathbf{W}] \right\}, \\ \text{Probabilities} \rightarrow \text{Regular}, \\ \text{BypassSearch} \rightarrow \text{True} \\ \left. \right] \\ \\ \left\{ \left\{ \mathbf{T} \rightarrow \{a_2, a_4\}, \right. \right. \\ \left. \left. \mathbf{W} \rightarrow \{a_3, a_4\}, \Omega \rightarrow \{a_1, a_2, a_3, a_4\} \right\}, \right. \\ \left. \left\{ a_1 \rightarrow \frac{6033}{8192}, a_2 \rightarrow \frac{1007}{8192}, a_3 \rightarrow \frac{1}{64}, a_4 \rightarrow \frac{1}{8} \right\} \right\} \end{array} \right]$$

TruthTable [MODEL4]

T	W	var	Pr
T	T	a_4	$\frac{1}{8}$
T	F	a_2	$\frac{1007}{8192}$
F	T	a_3	$\frac{1}{64}$
F	F	a_1	$\frac{6033}{8192}$

■ A PrSAT model showing that Sobel's (1)-(3) do *not* entail (5)

$$\text{MODEL5} = \text{PrSAT} \left[\begin{array}{l} \left\{ \text{Pr}[\mathbf{T}] < \frac{1}{2}, \right. \\ \quad \text{Pr}[\mathbf{T} \mid \mathbf{W}] > \frac{1}{2}, \\ \quad \text{Pr}[\mathbf{W} \mid \mathbf{T}] > \frac{1}{2}, \\ \quad \text{Pr}[\mathbf{T} \mid \mathbf{W}] - \text{Pr}[\mathbf{T} \mid \neg \mathbf{W}] < \text{Pr}[\neg \mathbf{W}] - \text{Pr}[\mathbf{W}] \\ \left. \right\}, \\ \text{Probabilities} \rightarrow \text{Regular}, \\ \text{BypassSearch} \rightarrow \text{True} \\ \left. \right] \end{array}$$

$$\left\{ \left\{ \mathbf{T} \rightarrow \{a_2, a_4\}, \right. \right. \\ \quad \left. \left. \mathbf{W} \rightarrow \{a_3, a_4\}, \Omega \rightarrow \{a_1, a_2, a_3, a_4\} \right\}, \right. \\ \left. \left\{ a_1 \rightarrow \frac{51}{64}, a_2 \rightarrow \frac{1}{16}, a_3 \rightarrow \frac{1}{16}, a_4 \rightarrow \frac{5}{64} \right\} \right\}$$

TruthTable [MODEL5]

T	W	var	Pr
T	T	a_4	$\frac{5}{64}$
T	F	a_2	$\frac{1}{16}$
F	T	a_3	$\frac{1}{16}$
F	F	a_1	$\frac{51}{64}$

- A PrSAT “Proof” that Sobel’s (1)-(3) do entail the disjunction (4) \vee (5)

```

PrSAT [
  { Pr [T] < 1/2,
    Pr [T | W] > 1/2,
    Pr [W | T] > 1/2,
    Not [Or [Pr [T | ¬ W] < Pr [W],
             Pr [T | W] - Pr [T | ¬ W] > Pr [¬ W] - Pr [W] ] ]
  },
  BypassSearch → True
]
{}

```