

Probabilistic Coherence from a Logical Point of View

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- Overview of the Talk
 - Foundation: Probabilistic Confirmation (c) from a Logical POV
 - * $c(h, e)$ as a “relevant” quantitative generalization of $\Box(e \supset h)$
 - * $c(h, e)$, so understood, is not $\Pr(e \supset h)$ or $\Pr(h|e)$, etc.
 - * $c(h, e)$ is something akin (ordinally) to the likelihood ratio
 - Defining Coherence (\mathcal{C}) in terms of “Mutual c -Confirmation”
 - * $\mathcal{C}(p, q)$ as a “mutual confirmation” generalization of $\Diamond(p \& q)$
 - * $\mathcal{C}(p, q)$, so understood, is not $\Pr(p \& q)$ or $\Pr(q|p)$, etc.
 - * Suggestion: $\mathcal{C}(p, q)$ as a function of $c(p, q)$ and $c(q, p)$, etc.
 - Confirmation as primitive, and coherence defined in terms of it
 - New definition of my \mathcal{C} measure (inspired by Moretti/Douven)
 - Some Subtleties/Objections (I’ll focus on “logical” ones)

From Confirmation to Coherence I

Confirmation (c)	Coherence (\mathcal{C})
Metatheoretic Concept: $\Box(e \supset h)$	Metatheoretic Concept: $\Diamond(p \& q)$
$\therefore e \models \sim h \Rightarrow$ maximal disconfirmation	$\therefore p \not\models \sim q \Rightarrow$ maximal incoherence
$\therefore e \models h [e \not\models \perp] \Rightarrow$ maximal confirmation	$\therefore p \models q \not\models \perp \Rightarrow$ maximal coherence
+ Dependence is confirmation	+ Dependence is coherence
– Dependence is disconfirmation	– Dependence is incoherence
Independence is neutrality	Independence is neutrality
$\Pr(e \supset h)$ won’t work	$\Pr(p \& q)$ won’t work
$\Pr(h e)$ won’t work, etc.	$\Pr(q p)$ won’t work, etc.
Most relevance measures won’t work	Most relevance measures won’t work

- In the confirmation case, only a small class of candidate c -measures will work.
- And, if \mathcal{C} is defined in terms of “mutual c ”, there are also few candidates.

From Confirmation to Coherence II

- Strategy: We will construct our \mathcal{C} measure using one of the proper c measures.
- We use a slight [new!] modification of Kemeny and Oppenheim’s c -measure F

$$F_M(h, e) =_{df} \begin{cases} \frac{\Pr_M(e|h) - \Pr_M(e|\sim h)}{\Pr_M(e|h) + \Pr_M(e|\sim h)} & \text{if } e \not\models h \text{ and } e \not\models \sim h. \\ 1 & \text{if } e \models h, \text{ and } e \not\models \perp. \\ -1 & \text{if } e \models \sim h. \end{cases}$$

- Let \mathcal{F} be the set containing the F values of all pairs of conjunctions of (thanks, Igor!) nonempty, disjoint subsets of the set of statements. And, \mathcal{C} is an average of \mathcal{F} . Note: F (hence \mathcal{C}) is *relativized* to a (regular) Pr-model M !
- \mathcal{F} is non-trivial to visualize! I haven’t analyzed the combinatorics of \mathcal{F} yet, but I have an algorithm for generating it. See my *MATHEMATICA*® notebook.
- I first proposed simply taking the straight average of \mathcal{F} , but other averages could be given (undoubtedly, some examples will suggest unequal weights).

Some Subtleties/Objections

- Individuation: The “information sets” (collections that \mathcal{C} measures) could be *multisets/sequences of propositions*, or *sets of statements* (tokens), etc., but not *sets of propositions*, unless we go anti-Stalnaker (which is controversial).
- Siebel: “if we are confronted with a pair of statements which cannot both be false together, Fitelson’s function assigns it a coherence value of at most 0.”
- True. But, this will be true for *any* Pr-relevance-based account (not just mine). If p and q can’t both be *false*, then they cannot be *positively* correlated! Here, correlation goes beyond a naïve generalization of the metatheoretic $\Diamond(p \& q)$.
- Moretti (and others): On your view, logically equivalent sets of statements can have different degrees of coherence. Yep. But, this also strikes me as correct. [To my mind, $\{p, q, r\}$ is more coherent than $\{p, q\}$, provided that $r \models p$.]
- Moretti: But, on your \mathcal{C} , adding \top to a coherent set can make it *incoherent*! This *was* true on my old \mathcal{C} . But, *not* on my new \mathcal{C} . See my *MATHEMATICA*® notebook. $\mathcal{C}(\mathbf{S})$ can be $< \mathcal{C}(\mathbf{S} \cup \{\top\})$, but this is an artifact of *averaging*.