

Lecture 3: Comparative Confidence

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- Kenny has written a paper [8] that explains how to relax the assumption of Opinionation in our framework.
- Our approach is equivalent to assigning (in)accurate judgments a *score* of $(-1) + 1$, and calculating the *total score* of \mathbf{B} (at w) as the *sum* of the scores of all $p \in \mathcal{A}$.
- Kenny's Generalizations: (a) allow scores of $-w$ and $+r$, where $w \geq r > 0$, and (b) allow S to *suspend on* p [$S(p)$], where all suspensions are given a *neutral* score of *zero*.
- This generalization of our framework leads to an elegant analogue of our central Theorem that (\mathcal{R}) entails (WADA).

Theorem. An agent S will avoid (strict) dominance in *total score* **if** their belief set \mathbf{B} can be represented as follows:


(\mathcal{R}) There exists a probability function $\text{Pr}(\cdot)$ such that, $\forall p \in \mathcal{A}$:


$$B(p) \text{ iff } \text{Pr}(p) > \frac{w}{r+w},$$

$$D(p) \text{ iff } \text{Pr}(p) < 1 - \frac{w}{r+w},$$

$$S(p) \text{ iff } \text{Pr}(p) \in \left[1 - \frac{w}{r+w}, \frac{w}{r+w}\right].$$

- We (along with Rachael Briggs and Fabrizio Cariani) [1] are investigating various applications of this new approach.
- One interesting application is to *judgment aggregation*. E.g.,
 - Majority rule aggregations of the judgments of a bunch of agents — each of whom satisfy (PV) — *need not* satisfy (PV).
- **Q:** does majority rule preserve *our* notion of coherence, viz., is (WADA) preserved by MR? **A:** yes (on simple, atomic + truth-functional agendas), but *not on all possible agendas*.
 - There are (not merely atomic + truth-functional) agendas A and sets of judges J ($|A| \geq 5$, $|J| \geq 5$) that (severally) satisfy (WADA), while their majority profile *violates* (WADA).
- *But*, if a set of judges is (severally) *consistent* [i.e., satisfy (PV)], then their majority profile *must* satisfy (WADA).

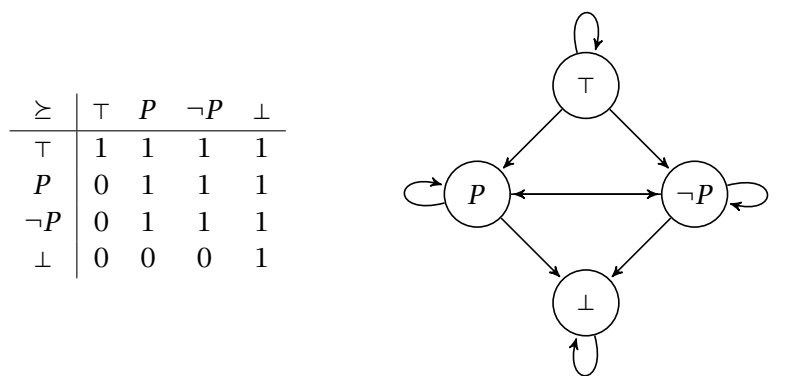
 **Recipe.** Wherever **B-consistency** runs into paradox, substitute *coherence* (in *our* sense), and see what happens.

- The contemporary literature focuses mainly on two types of *non-comparative* judgment: belief and credence. Not much attention is paid to *comparative* judgment (but see [16]).
 - It wasn't always thus. Keynes [21], de Finetti [3, 4] and Savage [24] all emphasized the importance (and perhaps even *fundamentality*) of comparative confidence.
 - *Comparative confidence* is a three-place relation between an agent S (at a time t) and a pair of propositions $\langle p, q \rangle$.
 - ' $p \succ q$ ' means ' S is strictly more confident in the truth of p than she is in the truth of q '. And, ' $p \sim q$ ' means ' S is equally confident in the truth of p and the truth of q '.
-  It's important to think of \succ and \sim as *autonomous* and *irreducibly comparative* — i.e., as a kind of comparative judgment *that may not reduce to anything non-comparative*.

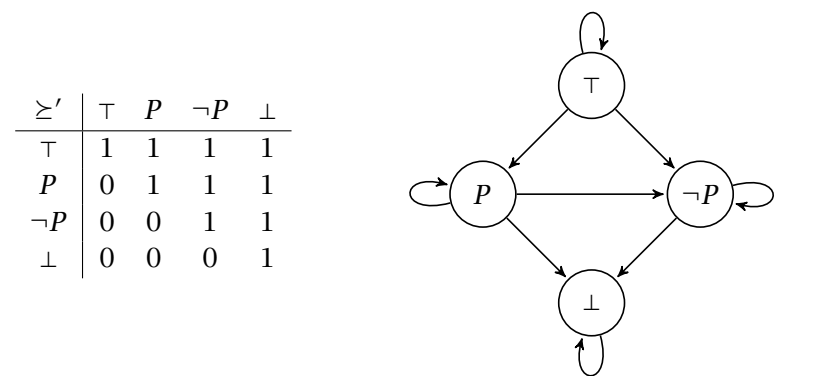
- There are good reasons to think that $p \geq q$ is not reducible to the credal comparison $b(p) \geq b(q)$, e.g., the dart case.

- First, we assume that S forms \succeq judgments regarding *all pairs* of propositions on some m -proposition agenda \mathcal{A} , drawn from an n -proposition Boolean algebra \mathcal{B}_n ($m \leq n$).
- Second, we assume that $>$ constitutes a *strict order* (on \mathcal{A}). That is, $>$ satisfies the following two ordering conditions.
 - Irreflexivity of $>$.** $p \not> p$.
 - Transitivity of $>$.** If $p > q$ and $q > r$, then $p > r$.
- Third, we assume that \sim is an *equivalence relation* on \mathcal{A} .
 - Reflexivity of \sim .** $p \sim p$.
 - Transitivity of \sim .** If $p \sim q$ and $q \sim r$, then $p \sim r$.
 - Symmetry of \sim .** If $p \sim q$, then $q \sim p$.
- Fourth, we assume our agents S are *logically omniscient*. (LO) S respects all logical equivalencies.
 - ☞ \therefore If p, q are logically equivalent, then S judges $p \sim q$. And, if S judges $p > q$, then p, q are *not* logically equivalent.

- Finally, we assume our agents S have *regular* \succeq -orderings.
 - Regularity.** If p is contingent, then $p > \perp$ and $\top > p$.
- This assumption threatens the application of the present framework to infinite underlying possibility spaces (especially, uncountable underlying possibility spaces). But, as we'll see below, it is crucial for many of our arguments.
- We can represent \succeq -relations on agendas \mathcal{A} via their 0/1 adjacency matrices A^\succeq , which are defined as follows:
 - $p_i > p_j$ iff $A_{ij}^\succeq = 1$ and $A_{ji}^\succeq = 0$.
 - $p_i \sim p_j$ iff $A_{ij}^\succeq = 1$ and $A_{ji}^\succeq = 1$.
- Toy example: let $\mathcal{A} = \mathcal{B}_4$ be the smallest sentential BA, with four propositions $\langle \top, P, \neg P, \perp \rangle$, for some contingent P . Specifically, interpret P as “a tossed coin lands heads.”
- The following \succeq relation seems natural, given this setup.



- The above figure shows the adjacency matrix and graphical representation of a relation (\succeq) on \mathcal{B}_4 . This relation \succeq is supported by S 's evidence E , if E says that the coin is *fair*.
- Consider an alternative relation (\succeq') on \mathcal{B}_4 , which agrees with \succeq on all judgments, *except for* $\neg P \succeq P$. That is, $P \succ' \neg P$; whereas, $P \sim \neg P$. [\succeq' is depicted on the next slide.]



- This alternative relation \succeq' on \mathcal{B}_4 is supported by S 's evidence E , if E says that the coin is *biased toward heads*.
- Intuitively, neither \succeq nor \succeq' should be deemed (formally) *incoherent*. After all, either could be supported by an agent's evidence. We'll return to evidential requirements for comparative confidence relations below. Meanwhile, Step 1.

Two Leftovers ○○	Background on \succeq ○○○○○	Coherence Requirements for \succeq ●○○○○○○○○○	Extras ○○○○○○○○○	Refs
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- **Step 1** involves articulating a precise sense in which an individual comparative confidence judgment $p \succeq q$ is *inaccurate* at w . Here, we follow Joyce’s [18, 19] *extensionality* assumption, which requires “inaccuracy” to *supervene on the truth-values of the propositions in \mathcal{A} at w* .
- ☞ An individual comparative confidence judgment $p \succeq q$ is *inaccurate at w* iff $p \succeq q$ entails that the ordering \succeq fails to rank all truths strictly above all falsehoods at w .¹
- On this conception, there are *two facts* about the inaccuracy of individual comparative confidence judgments $p \succeq q$.
 - Fact 1. If $q \ \& \ \neg p$ is true at w , then $p \succ q$ is inaccurate at w .
 - Fact 2. If $p \neq q$ is true at w , then $p \sim q$ is inaccurate at w .

¹One might be tempted by a weaker (and “more Joycean”) definition of inaccuracy, according to which $p \succeq q$ is inaccurate iff it *contradicts* the comparison $p \succeq_w q$ induced by the indicator function v_w . This weaker definition (which *also* deems $p \succ q$ inaccurate if $p \equiv q$ is true at w) is *untenable* for us. This will follow from our Fundamental Theorem, below.

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Two Leftovers ○○	Background on \succeq ○○○○○	Coherence Requirements for \succeq ●○○○○○○○○○	Extras ○○○○○○○○○	Refs
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- **Step 2** requires a *point-wise* inaccuracy measure $i(p \succeq q, w)$.
 - ☞ There are two kinds of inaccurate \succeq -judgments (Facts 1 and 2). *Intuitively, these two should kinds of inaccuracies should not receive equal i -scores.* Mistaken \succ judgments should receive *greater i -scores* than mistaken \sim judgments.
- *How much more inaccurate* than \sim mistakes are \succ mistakes? *Twice as inaccurate!* Suppose (by convention) that we assign an i -score of 1 to mistaken \sim judgments. We *must (!)* assign an i -score of 2 to mistaken \succ judgments.

$$i(p \succeq q, w) \stackrel{\text{def}}{=} \begin{cases} 2 & \text{if } q \ \& \ \neg p \text{ is true at } w, \text{ and } p \succ q, \\ 1 & \text{if } p \neq q \text{ is true at } w, \text{ and } p \sim q, \\ 0 & \text{otherwise.} \end{cases}$$
- \succeq 's *total inaccuracy* (on \mathcal{A} at w) is the *sum* of \succeq 's i -scores.

$$\mathcal{I}(\succeq, w) \stackrel{\text{def}}{=} \sum_{p, q \in \mathcal{A}} i(p \succeq q, w).$$

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Two Leftovers ○○	Background on \succeq ○○○○○	Coherence Requirements for \succeq ●○○○○○○○○○	Extras ○○○○○○○○○	Refs
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- **Step 3** involves the adoption of a *fundamental epistemic principle*. Here, we will follow Joyce and adopt:
 - Weak Accuracy-Dominance Avoidance (WADA).** \succeq should *not be weakly dominated* in inaccuracy (according to \mathcal{I}). More formally, there should *not* exist a \succeq' (on \mathcal{A}) such that
 - (i) $(\forall w) [\mathcal{I}(\succeq', w) \leq \mathcal{I}(\succeq, w)]$, and
 - (ii) $(\exists w) [\mathcal{I}(\succeq', w) < \mathcal{I}(\succeq, w)]$.
- Recall our toy relations \succeq and \succeq' over \mathcal{B}_4 . Neither of these relations should be *ruled-out as incoherent*, as each *could be* supported by *some* body of evidence [19, pp. 282–3].
- ☞ **Theorem.** *Neither \succeq nor \succeq' is weakly dominated in \mathcal{I} -inaccuracy — by any binary relation on \mathcal{B}_4 .*
- This result is a corollary of our Fundamental Theorem, which will also explain why we were *forced* to assign an inaccuracy score of *exactly 2* to inaccurate \succ judgments.
- More on that later. Meanwhile, a historical interlude.

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Two Leftovers ○○	Background on \succeq ○○○○○	Coherence Requirements for \succeq ●○○○○○○○○○	Extras ○○○○○○○○○	Refs
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- Various coherence requirements for \succeq have been discussed [15, 2, 26]. We’ll focus on a *particular family* of these.
- We begin with *the fundamental requirement* (\mathbb{C}), which has (near) universal acceptance. We will state (\mathbb{C}) in two ways: axiomatically, and in terms of numerical representability.
 - (\mathbb{C}) S 's \succeq -relation (assumed to be a total preorder on \mathcal{B}_n) should satisfy the following two axiomatic constraints:
 - (A₁) $\top \succ \perp$.
 - (A₂) For all $p, q \in \mathcal{B}_n$, if p entails q then $q \succeq p$.
- A *plausibility measure* (a.k.a., a *capacity*) on a Boolean algebra \mathcal{B}_n is real-valued function $\text{Pl} : \mathcal{B}_n \rightarrow [0, 1]$ which satisfies the following three conditions [15, p. 51]:
 - (Pl₁) $\text{Pl}(\perp) = 0$.
 - (Pl₂) $\text{Pl}(\top) = 1$.
 - (Pl₃) For all $p, q \in \mathcal{B}_n$, if p entails q then $\text{Pl}(q) \geq \text{Pl}(p)$.

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Two Leftovers	Background on \succeq	Coherence Requirements for \succeq	Extras	Refs
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<ul style="list-style-type: none"> Two kinds of representability of \succeq, by a real-valued f. <ul style="list-style-type: none"> \succeq is <i>fully</i> represented by $f \cong$ for all $p, q \in \mathcal{B}_n$ $p \succeq q \iff f(p) \geq f(q).$ \succeq is <i>partially</i> represented by $f \cong$ for all $p, q \in \mathcal{B}_n$ $p \succ q \implies f(p) > f(q).$ Now, (C) can be expressed equivalently, as follows: <ul style="list-style-type: none"> (C) S's \succeq-relation (assumed to be a total preorder on \mathcal{B}_n) should be <i>fully representable by some plausibility measure</i>. <p> Theorem 1. (WADA) entails (C). [See Extras for a proof.]</p> <ul style="list-style-type: none"> There are several other coherence requirements for \succeq that can be expressed both axiomatically, and in terms of numerical representability by some real-valued f. We'll state these, and say whether or not they follow from (WADA). The next requirements involve <i>belief functions</i>. 				
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Two Leftovers	Background on \succeq	Coherence Requirements for \succeq	Extras	Refs
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<ul style="list-style-type: none"> A <i>mass function</i> on a Boolean algebra \mathcal{B}_n is a function $m : \mathcal{B}_n \rightarrow [0, 1]$ that satisfies the following two conditions: <ul style="list-style-type: none"> (M₁) $m(\perp) = 0$. (M₂) $\sum_{p \in \mathcal{B}_n} m(p) = 1$. A <i>belief function</i> $\text{Bel} : \mathcal{B}_n \rightarrow [0, 1]$ is generated by an underlying mass function m on \mathcal{B}_n in the following way: $\text{Bel}_m(p) \stackrel{\text{def}}{=} \sum_{\substack{q \in \mathcal{B}_n \\ q \text{ entails } p}} m(q).$ Now, consider the following coherence requirement: <ul style="list-style-type: none"> (C₀) S's \succeq-relation (assumed to be a total preorder on \mathcal{B}_n) should be <i>partially</i> representable by some belief function. A total preorder \succeq satisfies (C₀) iff \succeq satisfies (A₂) [26]. So, Theorem 1 has a Corollary: ["Thm 2"] (WADA) entails (C₀). What about <i>full</i> representability of a belief function? To wit: 				
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Two Leftovers	Background on \succeq	Coherence Requirements for \succeq	Extras	Refs
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<ul style="list-style-type: none"> (C₁) S's \succeq-relation (assumed to be a total preorder on \mathcal{B}_n) should be <i>fully</i> representable by a belief function. As it turns out [26], a relation \succeq is <i>fully</i> representable by some belief function if and only if \succeq satisfies (A₁), (A₂), and <ul style="list-style-type: none"> (A₃) If p entails q and $\langle q, r \rangle$ are mutually exclusive, then: $q \succ p \implies q \vee r \succ p \vee r.$ (WADA) also entails (A₃). That is, we have the following: <p>Theorem 3. (WADA) entails (C₁). [See Extras.]</p> Moving beyond (C₁) takes us into <i>comparative probability</i>. A t.p. \succeq is a <i>comparative probability</i> iff \succeq satisfies (A₁), (A₂), & <ul style="list-style-type: none"> (A₅) If $\langle p, q \rangle$ and $\langle p, r \rangle$ are mutually exclusive, then: $q \succeq r \iff p \vee q \succeq p \vee r$ (C₂) S's \succeq-relation (assumed to be a total preorder on \mathcal{B}_n) should be a <i>comparative probability</i> relation. 				
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Two Leftovers	Background on \succeq	Coherence Requirements for \succeq	Extras	Refs
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<p>Theorem 4. (WADA) does <i>not</i> entail (C₂). [See Extras.]</p> <ul style="list-style-type: none"> The following axiomatic constraint is a weakening of (A₅). <ul style="list-style-type: none"> (A₅[*]) If $\langle p, q \rangle$ and $\langle p, r \rangle$ are mutually exclusive, then: $q \succ r \implies p \vee r \not\succeq p \vee q$ And, the following coherence requirement is a (corresponding) weakening of coherence requirement (C₂). <ul style="list-style-type: none"> (C₂[*]) \succeq should (be a total preorder and) satisfy (A₁), (A₂) and (A₅[*]). Theorem 5. (WADA) does <i>not</i> entail (C₂[*]). [See Extras.] Our final pair of coherence requirements for \succeq involve representability by some <i>probability</i> function. I'm sure everyone knows what a Pr-function is, but... Probability functions are special kinds of belief functions (just as belief functions were special kinds of Pl-measures). 				
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- A *probability mass function* is a function m which maps *states* of \mathcal{B}_n to $[0, 1]$, and which satisfies these two axioms.

$$(N_1) \quad m(\perp) = 0.$$

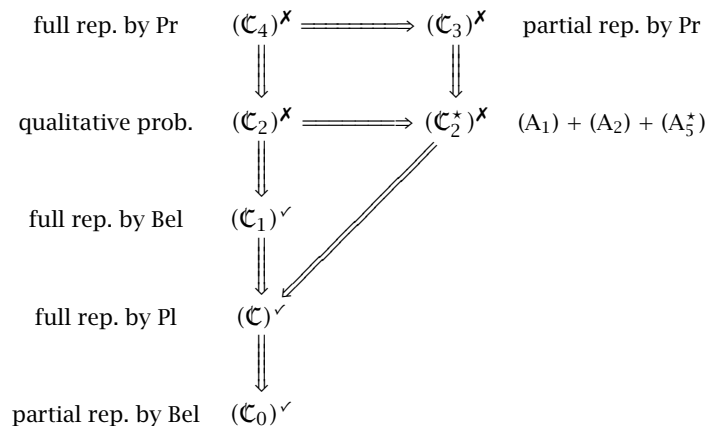
$$(N_2) \quad \sum_{s \in \mathcal{B}_n} m(s) = 1.$$

- A *probability function* $\text{Pr} : \mathcal{B}_n \rightarrow [0, 1]$ is generated by an underlying probability mass function m in the following way

$$\text{Pr}_m(p) \stackrel{\text{def}}{=} \sum_{\substack{s \in \mathcal{B}_n \\ s \text{ entails } p}} m(s).$$

- That brings us to our final pair of requirements for \succeq .
 - (C_3) \succeq should be *partially* representable by some Pr-function.
 - (C_4) \succeq should be *fully* representable by some Pr-function.
- de Finetti [3, 4] famously conjectured that (C_2) entails (C_4) . But, Kraft *et. al.* [22] showed that $(C_2) \not\Rightarrow (C_3)$. [See Extras.]

- We have the following logical relations between the C 's.



- If a requirement follows from (WADA), it gets a “ \checkmark ”. If a requirement does *not* follow from (WADA), it gets an “ X ”.
- We conclude with our final (and most important) Fundamental Theorem(s). [See Extras for proofs.]

- We assume that “numerical probabilities reflect evidence”, *i.e.*, we adopt the following *evidential requirement*.

(R) \succeq is representable by some *regular* probability function.

Fundamental Theorem. If a comparative confidence relation \succeq satisfies (R) , then \succeq satisfies (WADA).

- The proof of our Fundamental Theorem (see Extras) reveals that $\mathcal{I}(\succeq, w)$ is *evidentially proper*, in this sense [13].

Definition (Evidential Propriety). Suppose a judgment set J of type \mathcal{J} is supported by the evidence. That is, suppose there exists some evidential probability function $\text{Pr}(\cdot)$ which represents J (in the appropriate sense of “represents” for judgment sets of type \mathcal{J}). If this is sufficient to ensure that J minimizes expected inaccuracy (relative to Pr), according to the measure of inaccuracy $\mathcal{I}(J, w)$, then we will say that the measure \mathcal{I} is **evidentially proper**.

Note: the decision to weight \succ -mistakes *twice as heavily* as \sim -mistakes is *forced* by evidential propriety (see Extras).

Theorem 1. (WADA) entails (C) , *viz.*, $(\text{WADA}) \Rightarrow (A_1) \ \& \ (A_2)$.

Proof.

Suppose \succeq violates (A_1) . Because \succeq is total, this means \succeq is such that $\perp \succeq \top$. Consider the relation \succeq' which agrees with \succeq on all comparisons outside the $\langle \perp, \top \rangle$ -fragment, but which is such that $\top \succ' \perp$. We have: $(\forall w) [\mathcal{I}(\top \succ' \perp, w) = 0 < 1 \leq \mathcal{I}(\perp \succeq \top, w)]$. \square

Suppose \succeq violates (A_2) . Because \succeq is total, this means there is a pair of propositions p and q in \mathcal{A} such that (a) p entails q but (b) $p \succ q$. Consider the relation \succeq' which agrees with \succeq outside of the $\langle p, q \rangle$ -fragment, but which is such that $q \succ' p$. The table on the next slide depicts the $\langle p, q \rangle$ -fragments of the relations \succeq and \succeq' in the three salient possible worlds w_1 - w_3 not ruled out by (a) $p \models q$. By (b) & (LO), p and q are not logically equivalent. So, world w_2 is a live possibility, and \succeq' weakly \mathcal{I} -dominates \succeq . \square

Two Leftovers: ○○ Background on \succeq : ○○○○ Coherence Requirements for \succeq : ○○○○○○○○○ Extras: ●○○○○○○○ Refs: ○○

w_i	p	q	\succeq	\succeq'	$\mathcal{I}(\succeq, w_i)$	$\mathcal{I}(\succeq', w_i)$
w_1	T	T	$p \succ q$	$q \succ' p$	0	0
	T	F				
w_2	F	T	$p \succ q$	$q \succ' p$	2	0
w_3	F	F	$p \succ q$	$q \succ' p$	0	0

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Two Leftovers: ○○ Background on \succeq : ○○○○ Coherence Requirements for \succeq : ○○○○○○○○○ Extras: ●○○○○○○○ Refs: ○○

Theorem 3. (WADA) entails (\mathcal{C}_1) .

Proof.
 Having already proved Theorem 1, we just need to show that (WADA) entails (A_3) . Suppose \succeq violates (A_3) . Because \succeq is total, this means there must exist $p, q, r \in \mathcal{A}$ such that (a) $p \models q$, (b) $\langle q, r \rangle$ are mutually exclusive, (c) $q \succ p$, but (d) $p \vee r \succeq q \vee r$. Let \succeq' agree with \succeq on every judgment, *except* (d). That is, let \succeq' be such that (e) $q \succ' p$ and (f) $q \vee r \succ' p \vee r$. There are only four worlds (or $\langle p, q, r \rangle$ state descriptions) compatible with the precondition of (A_3) . These are the following (state descriptions).

$$w_1 = p \& q \& \neg r \quad w_2 = \neg p \& q \& \neg r$$

$$w_3 = \neg p \& \neg q \& r \quad w_4 = \neg p \& \neg q \& \neg r$$

By (c) & (LO), p and q are not logically equivalent. As a result, world w_2 is a live possibility. Moreover, (f) will *not* be inaccurate in *any* of these four worlds. But, (d) *must be inaccurate in world w_2* . This suffices to show that \succeq' weakly \mathcal{I} -dominates \succeq . \square

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Two Leftovers: ○○ Background on \succeq : ○○○○ Coherence Requirements for \succeq : ○○○○○○○○○ Extras: ○○○●○○○○○ Refs: ○○

Theorem 4. (WADA) does *not* entail (\mathcal{C}_2) .

Proof.
 Having already proved Theorem 1, we just need to show that (WADA) does *not* entail (A_5) . Suppose (a) $\langle p, q \rangle$ and $\langle p, r \rangle$ are mutually exclusive, (b) $q \succ r$, and (c) $p \vee r \succ p \vee q$. It can be shown (by exhaustive search) that *there is no binary relation \succeq' on the agenda $\langle p, q, r \rangle$ such that (i) \succeq' agrees with \succeq on all judgments *except* (b) and (c), and (ii) \succeq' weakly \mathcal{I} -dominates \succeq* . There are only four alternative judgment sets that need to be compared with $\{(b), (c)\}$, in terms of their \mathcal{I} -values across the five possible worlds (w_1 - w_5) compatible with the precondition of (A_5) : (1) $\{q \sim r, p \vee r \succ p \vee q\}$, (2) $\{r \succ q, p \vee r \succ p \vee q\}$, (3) $\{q \succ r, p \vee r \sim p \vee q\}$, and (4) $\{q \sim r, p \vee r \sim p \vee q\}$. It is easy to verify that none of these alternative judgment sets weakly \mathcal{I} -dominates the set $\{(b), (c)\}$, across the five salient possible worlds. Note: this argument actually establishes the *stronger* claim (**Theorem 5**) that (WADA) does *not* entail $(A_5^*)/(\mathcal{C}_2^*)$. \square

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Two Leftovers: ○○ Background on \succeq : ○○○○ Coherence Requirements for \succeq : ○○○○○○○○○ Extras: ○○○●○○○○○ Refs: ○○

Fundamental Theorem. If a comparative confidence relation \succeq satisfies (\mathcal{R}) , then \succeq satisfies (WADA). That is, $(\mathcal{R}) \Rightarrow$ (WADA).

Proof.
 Suppose $\text{Pr}(\cdot)$ fully represents \succeq . Consider the expected \mathcal{I} -inaccuracy, as calculated by $\text{Pr}(\cdot)$, of \succeq : $\mathbb{E}\mathcal{I}_{\text{Pr}}^{\succeq} \stackrel{\text{def}}{=} \sum_w \text{Pr}(w) \cdot \mathcal{I}(\succeq, w)$. Since $\mathcal{I}(\succeq, w)$ is a sum of the $i(p \succeq q, w)$ for each $\langle p, q \rangle \in \mathcal{A}$, and since \mathbb{E} is linear:

$$\mathbb{E}\mathcal{I}_{\text{Pr}}^{\succeq} = \sum_{p, q \in \mathcal{A}} \mathbb{E}_{\text{Pr}} i(p \succeq q, w)$$

- Suppose $\text{Pr}(p) > \text{Pr}(q)$. Then we have:
 $\mathbb{E}_{\text{Pr}} i(p \succ q, w) = 2 \cdot \text{Pr}(q \& \neg p) < \mathbb{E}_{\text{Pr}} i(p \sim q, w) = \text{Pr}(p \neq q)$, and
 $\mathbb{E}_{\text{Pr}} i(p \succ q, w) = 2 \cdot \text{Pr}(q \& \neg p) < \mathbb{E}_{\text{Pr}} i(q \succ p, w) = 2 \cdot \text{Pr}(p \& \neg q)$.
- Suppose $\text{Pr}(p) = \text{Pr}(q)$. Then we have:
 $\mathbb{E}_{\text{Pr}} i(p \sim q, w) = \text{Pr}(p \neq q) = \mathbb{E}_{\text{Pr}} i(p \succ q, w) = 2 \cdot \text{Pr}(q \& \neg p)$.

As a result, if \succeq is fully representable by *any* $\text{Pr}(\cdot)$, then \succeq cannot be *strictly \mathcal{I} -dominated*, i.e., $(\mathcal{C}_4) \Rightarrow$ (SADA). Moreover, if we assume $\text{Pr}(\cdot)$ to be *regular*, then \succeq must satisfy (WADA) [13]. $\therefore (\mathcal{R}) \Rightarrow$ (WADA). \square

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Theorem. $a := 2; b := 0$ is the only assignment to a, b that ensures the following definition of i is *evidentially proper*.

$$i(p \succeq q, w) \stackrel{\text{def}}{=} \begin{cases} a & \text{if } q \ \& \ \neg p \text{ is true in } w, \text{ and } p \succ q, \\ b & \text{if } q \equiv p \text{ is true in } w, \text{ and } p \succ q, \\ 1 & \text{if } p \not\equiv q \text{ is true in } w, \text{ and } p \sim q, \\ 0 & \text{otherwise.} \end{cases}$$

Let $m_4 = \Pr(p \ \& \ q)$, $m_3 = \Pr(\neg p \ \& \ q)$, and $m_2 = \Pr(p \ \& \ \neg q)$. Then, the propriety of i is equivalent to the following (universal) claim. And, the only assignment that makes this (universal) claim true is $a := 2; b := 0$.

$$m_2 + m_4 > m_3 + m_4 \Rightarrow \left(\begin{array}{c} a \cdot m_3 + b \cdot (1 - (m_2 + m_3)) \leq a \cdot m_2 + b \cdot (1 - (m_2 + m_3)) \\ \& \\ a \cdot m_3 + b \cdot (1 - (m_2 + m_3)) \leq m_2 + m_3 \end{array} \right)$$

&

$$m_2 + m_4 = m_3 + m_4 \Rightarrow \left(\begin{array}{c} m_2 + m_3 \leq a \cdot m_2 + b \cdot (1 - (m_2 + m_3)) \\ \& \\ m_2 + m_3 \leq a \cdot m_3 + b \cdot (1 - (m_2 + m_3)) \end{array} \right)$$

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- Our ordering presuppositions (Totality & Transitivity) are not universally accepted as rational requirements [14, 12, 23].
- In our book [13], we analyze both of the ordering presuppositions in more detail. Specifically, we show that:
 - (1) Totality does not follow from weak accuracy dominance avoidance. That is, (WADA) does not entail Totality.
 - (2) Transitivity does not follow from weak accuracy dominance avoidance. That is, (WADA) does not entail Transitivity.
- These two negative results [especially (1)] are probably not very surprising. But, it is somewhat interesting that *none of the three instances of Transitivity is entailed by (WADA)*.
 - Transitivity₁.** If $p \succ q$ and $q \succ r$, then $r \succ p$.
 - Transitivity₂.** If $p \succ q$ and $q \sim r$, then $r \succ p$.
 - Transitivity₃.** If $p \sim q$ and $q \sim r$, then $p \sim r$.
- The first instance of Transitivity is the *least* controversial of the three. And, the last (transitivity of \sim) is the *most* [23].

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- There are two, weaker \mathcal{I} -dominance requirements that we discuss in the book [13]. These are as follows.

Strict Accuracy-Dominance Avoidance (SADA). \succeq should *not be strictly dominated* in inaccuracy (according to \mathcal{I}). More formally, there should *not* exist a \succeq' (on \mathcal{A}) such that

$$(\forall w) [\mathcal{I}(\succeq', w) < \mathcal{I}(\succeq, w)].$$
- Of course, (SADA) is *strictly weaker* than (WADA). And, here is a requirement that is *even weaker* than (SADA).
- Let $\mathbf{M}(\succeq, w) \stackrel{\text{def}}{=} \text{the set of } \succeq\text{'s in inaccurate judgments at } w$.

Strong Strict Accuracy-Dominance Avoidance (SSADA). There should *not* exist a \succeq' on \mathcal{A} such that:

$$(\forall w) [\mathbf{M}(\succeq', w) \subset \mathbf{M}(\succeq, w)].$$
- Some of our (WADA) results *also go through for* (SADA) and/or (SSADA). Finally, we give a complete, “big picture” of all the logical relations among all the requirements.

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Two Leftovers ○○ Background on \succeq ○○○○ Coherence Requirements for \succeq ○○○○○○○○○ Extras ○○○○○●○○ Refs

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Two Leftovers ○○	Background on \geq ○○○○	Coherence Requirements for \geq ○○○○○○○○○○	Extras ○○○○○○○○	Refs	
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Two Leftovers ○○	Background on \geq ○○○○	Coherence Requirements for \geq ○○○○○○○○○○	Extras ○○○○○○○○	Refs	
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