

SOME RECENT RESULTS IN ALGEBRA & LOGICAL CALCULI  
OBTAINED USING AUTOMATED REASONING

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<sup>a</sup>Results reported here are either the result of joint work, or the work of others associated with AR @ MCS @ ANL: Larry Wos, Bill McCune, Ken Kunen, Steve Winker, Bob Veroff, Ken Harris, Zac Ernst, John Slaney, Ted Ulrich, Bob Meyer, R. Padmanabhan *et al.*

## Overview of Presentation

- Simple Equational Bases for Boolean Algebra
  - In the connectives + (disjunction) and  $n$  (negation)
  - In the connective | (Sheffer Stroke)
  - Pointers to recent work in other algebraic systems
- Simple Axiomatizations of Classical Propositional Logics
  - Classical Propositional Logic (in two sets of connectives)
  - Equational Fragment of Classical Propositional Logic
- Simple Axiomatizations of Non-Classical Propositional Logics
  - Strict Implicational fragments of modal logics S4, S5
  - Implicational fragments of relevance logics R, RM, E,  $\mathcal{L}$
- Time permitting: discussion of methods

## Equational Bases for BA in + and $n$ I

- In 1933, E.V. Huntington presented the following 3-basis for BA [10, 9]:

(Commutativity+)  $x + y = y + x$

(Associativity+)  $(x + y) + z = x + (y + z)$

(Huntington)  $n(n(x) + y) + n(n(x) + n(y)) = x$

- BA is usually presented in terms of +, ·,  $n$ , 0, 1. From Huntington's basis, 0, 1, and ·, with appropriate properties, can be established (*easy* for OTTER [20]).
- Shortly thereafter, Herbert Robbins asked whether the Huntington equation can be replaced with the following equation (which is shorter by one “ $n$ ”):

(Robbins)  $n(n(x + y) + n(x + n(y))) = x$

- The Robbins problem remained open for over 63 years, and attracted the attention of various people, including Tarski, and others [8], [2].

## Equational Bases for BA in + and $n$ II

- In 1979, Steve Winker, a student visiting Argonne, learned of the Robbins problem from Joel Berman. He and Larry Wos began to attack the problem.
- Larry Wos suggested looking for properties that force Robbins algebras to be Boolean. Winker [52] ingeniously found several such conditions (both “hand” and automated reasoning), including the following two relatively weak ones:
  1.  $\exists c \exists d (c + d = c)$
  2.  $\exists c \exists d (n(c + d) = n(c))$
- In 1996, Bill McCune [22] used an Argonne TP (EQP [21], a cousin of OTTER [20]) to prove that all Robbins algebras satisfy Winker's (2), above.
- This solved the long-standing Robbins problem. But, the machine proof of Winker's condition was not very easy for a human to follow or understand.
- Since McCune's discovery, several people (including myself [7]) have tried, in various ways, to make the EQP (and OTTER) proofs easier to digest [3].

### Equational Bases for BA in + and $n$ III

- It was thought that the (32-symbol) Robbins basis for BA was the simplest known, until I dug-up the following 23-symbol 2-basis for Boolean algebra reported (without proof) by Carew Meredith in 1968 [30, p. 228]:

$$\text{(Meredith}_1) \quad n(n(x) + y) + x = x$$

$$\text{(Meredith}_2) \quad n(n(x) + y) + (z + y) = y + (z + x)$$

- In 1966, Tarski [45] reported that BA does have *single* +,  $n$  axioms. Building on Tarski's work, Padmanabhan and Quackenbush [33] gave a method for constructing such axioms. But, their method yields *long* single axioms [23].
- Recently, Bill McCune [26] discovered a 22-symbol single axiom for BA:

$$\text{(DN}_1) \quad n(n(n(x + y) + z) + n(x + n(n(z) + n(z + u)))) = z$$

- OPEN: do shorter single axioms (or bases!) for BA (in + and  $n$ ) exist?

### Equational Bases for BA in Sheffer's | I

- In 1913, Sheffer [42] gave the following 3-basis for Boolean algebra in terms of a single binary connective | (Sheffer's | is just NAND:  $x|y = n(x) + n(y)$ ).

$$\text{(Sheffer}_1) \quad (x|x)|(x|x) = x$$

$$\text{(Sheffer}_2) \quad x|(y|(y|y)) = x|x$$

$$\text{(Sheffer}_3) \quad (x|(y|z)|(x|(y|z))) = ((y|y)|x)|((z|z)|x)$$

- Meredith [28] (again, in obscurity, and rediscovered by me) simplified matters in 1969 by presenting the following (23-symbol) 2-basis for the same theory.

$$\text{(Meredith}_3) \quad (x|x)|(y|x) = x$$

$$\text{(Meredith}_4) \quad x|(y|(x|z)) = ((z|y)|y)|x$$

- Recently, Bob Veroff [50] established the following (17-symbol) 2-basis:

$$\text{(Commutativity |)} \quad x|y = y|x$$

$$\text{(Veroff}_{26a}) \quad (x|y)|(x|(y|z)) = x$$

### Equational Bases for BA in Sheffer's | II

- The work of Tarski [45] and Padmanabhan & Quackenbush [33] also implies the existence of single axioms for BA in the Sheffer Stroke. But, as before, the |-single axioms generated by the methods of [33] are quite long [23].

- Recently, [53, 26] the following 15-symbol 1-bases were discovered.

$$\text{(Sh}_1) \quad (x|((y|x)|x))|(y|(z|x)) = y$$

$$\text{(Sh}_2) \quad ((y|(x|y))|y)|(x|(z|y)) = x$$

- It is also known [53, 26] that these |-single axioms are the *shortest possible*.
- Elegant axioms for groups [13, 19], lattices [25], loops [14, 15, 12], and other algebraic structures [24] have been discovered by the extended Argonne team.

### Sheffer Stroke Single Axioms for Sentential Logic I

- In 1917, Nicod [32] showed<sup>a</sup> that the following 23-symbol formula (in Polish notation) is a single axiom for classical sentential logic ( $D$  is interpreted semantically as NAND, *i.e.*, the Sheffer stroke):

$$\text{(N)} \quad DDpDqrDDtDttDDsqDDpsDps$$

- The only rule of inference for Nicod's single axiom system is the following, somewhat odd, detachment rule for  $D$ :

$$\text{(D-Rule)} \quad \text{From } DpDqr \text{ and } p, \text{ infer } r.$$

- Łukasiewicz [17, pp. 179–196] later showed that the following *substitution instance* ( $t/s$ ) of Nicod's axiom (N) would suffice:

$$\text{(\mathcal{L}_1)} \quad DDpDqrDDsDssDDsqDDpsDps$$

<sup>a</sup>Actually, Nicod's original proofs are erroneous (as noted by Łukasiewicz in [17]). See Scharle's [41] for a rigorous proof of the completeness of Nicod's system.

### Sheffer Stroke Single Axioms for Sentential Logic II

- Łukasiewicz's student Mordchaj Wajsberg [51, pp. 37–39] later discovered the following *organic*<sup>a</sup> 23-symbol single axiom for  $D$ :

$$(W) \quad DDpDqrDDDsrDDpsDpsDpDpq$$

- Łukasiewicz later discovered another 23-symbol organic axiom:

$$(\mathbb{L}_2) \quad DDpDqrDDpDrpDDsqDDpsDps$$

- Ken Harris and I have recently discovered many new 23-symbol single axioms, some of which are organic and have only 4 variables, e.g.,

$$(HF_1) \quad DDpDqrDDpDqrDDsrDDrsDps$$

- We have also shown that 23 symbol axioms are the shortest possible.

<sup>a</sup>A single axiom is *organic* if it contains no tautologous subformulae. (N) and (L) are *non-organic*, because they contain tautologous subformulae of the form  $Dx Dxx$ .

### Single $C$ - $O$ Axioms for Classical Sentential Logic

- Meredith [27, 29] reports two 19-symbol single axioms for classical sentential logic (using only the rule of condensed detachment, or *modus ponens* for  $C$ ) in terms of implication  $C$  and the constant  $O$  (semantically,  $O$  is “The False”):

$$CCCCCpqCrOstCCtpCrp$$

$$CCCpqCCOrsCCspCtCup$$

- Meredith [27, page 156] claims to have “almost completed a proof that no single axiom of  $(C, O)$  can contain less than 19 letters.” As far as we know, no such proof was ever completed (that is, until now...).
- We have performed an exhaustive search/elimination of all  $(C, O)$  theorems with fewer than 19 symbols. We have proven Meredith's conjecture: *no single axiom of classical PL in  $(C, O)$  can contain less than 19 letters.*<sup>a</sup>

<sup>a</sup>The elimination of some  $(C, O)$  candidates relied on matrices generated using *stochastic local search* techniques (as described by Ted Ulrich in his [47, 49] and by Cipra [2]). Stochastic local search is very powerful in the context of implicational logics. It has led to *many* useful (small) models.

### Single Axioms for The Equivalential Fragment of Classical Sentential Logic

- In 1933, Łukasiewicz [46, 250–277] showed (*lots* of hand calculations!) that (with MP for  $E$  as the sole rule) the shortest single axioms for the equivalential ( $E$ ) fragment of classical propositional logic contain 11 symbols. He found 2 such axioms.
- In the 1950's, Meredith [29] discovered seven more 11-symbol single axioms for  $E$ .
- John Kalman [11], and his student J. Peterson [36, 37], did extensive work on the problem in the 1970's. They found one more 11-symbol single axiom, and they eliminated all but 7 of the remaining 640 11-symbol candidate single axioms.
- In 1977–1979, Wos, Winker, *et al* (all at Argonne) [55] worked on the remaining 7 candidates. They ruled-out all but three, and showed that two of these three were single axioms. This left the following (*and last*) remaining 11-symbol candidate:

$$(XCB) \quad EpEEEpqErqr$$

- About a year ago, we (Wos, Dolph Ulrich, Fitelson [54]) proved that XCB is a single axiom for the equivalential calculus. The proof contains substitution instances with over 2000 symbols. This completes a 70-year study initiated by Łukasiewicz.

### New Bases for $C5$

- In their classic paper [16], Lemmon, Meredith, Meredith, Prior, and Thomas present several axiomatizations (assuming only the rule of condensed detachment, or *modus ponens* for  $C$ ) of the system  $C5$ , which is the strict-implicational fragment of the modal logic  $S5$ .
- Bases for  $C5$  containing 4, 3, 2, and a single axiom are presented in [16]. The following 2-basis is the shortest of these bases. It contains 20 symbols, 5-variables, and 9 occurrences of the connective  $C$ .

$$Cpp$$

$$CCCCpqrqCCqsCtCps$$

- The following 21-symbol (6-variable, 10- $C$ ) single axiom (due to C.A. Meredith) for  $C5$  is also reported in [16]:

$$CCCCCtppqCrCsCCspCuCrp$$

### New Bases for C5 (Cont'd)

- We (Ernst, Fitelson, Harris, Vos) searched both for new (hopefully, shorter than previously known) single axioms for C5 and for new 2-bases for C5.
- We discovered the following new 2-basis for C5, which is shorter than any previously known basis (indeed, it is as short as *any possible* basis — see below). It has 18 symbols, 4 variables, and 8 occurrences of  $C$ :

$$C\ p\ p \\ C\ C\ p\ q\ C\ C\ C\ C\ q\ r\ s\ r\ C\ p\ r$$

- Moreover, we discovered the following new 21-symbol (6-variable, 10- $C$ ) single axiom for C5 (as well as 5 others, not given here):

$$C\ C\ C\ C\ p\ q\ r\ C\ C\ s\ s\ q\ C\ C\ q\ t\ C\ u\ C\ p\ t$$

- *No formula with fewer than 21 symbols is a single axiom for C5. And, no basis for C5 whatsoever has fewer than 18 symbols.* Results to appear in [5].

### New Bases for C4

- C4 is the strict-implicational fragment of the modal logic S4 (and several other modal logics in the neighborhood of S4 — see Ulrich's [48]).
- As far as we know, the shortest known basis for C4 is due to Ulrich (see Ulrich's [48]), and is the following 25-symbol, 11- $C$ , 3-axiom basis:

$$C\ p\ p\ p \qquad C\ C\ p\ q\ C\ r\ C\ p\ q \qquad C\ C\ p\ C\ q\ r\ C\ C\ p\ q\ C\ p\ r$$

- Anderson & Belnap [1, p. 89] state the finding of a (short) single axiom for C4 as an open problem (as far as we know, this has *remained* open). The following is a 21-symbol (6-variable, 10- $C$ ) single axiom for C4:

$$C\ C\ p\ C\ C\ q\ C\ r\ r\ C\ p\ s\ C\ C\ s\ t\ C\ u\ C\ p\ t$$

- We have also the following 20-symbol 2-basis for C4:

$$C\ p\ C\ q\ q \qquad C\ C\ p\ C\ q\ r\ C\ C\ p\ q\ C\ s\ C\ p\ r$$

- *No formula with fewer than 21 symbols is a single axiom for C4. And, no basis for C4 whatsoever has fewer than 20 symbols.* Results to appear in [5].

### New Bases for $RM_{\rightarrow}$

- The “classical” relevance logic R-Mingle (RM) was first carefully studied by Dunn in the late 60's (e.g., in [4]). Interestingly, the implicational fragment of R-Mingle ( $RM_{\rightarrow}$ ) has an older history.
- $RM_{\rightarrow}$  was studied (albeit, unwittingly!) by Sobociński in the early 50's. Sobociński [43] discusses a two-designated-value-variant of Łukasiewicz's three-valued implication-negation logic (I'll call Sobociński's logic  $S_{\rightarrow}$ ). Sobociński leaves the axiomatization of  $S_{\rightarrow}$  as an open problem.
- Rose [39, 40] solved Sobociński's open problem, but his axiomatizations of  $S_{\rightarrow}$  are very complicated and highly redundant (see Parks' [34]).
- Meyer & Parks [31, 35] report: (i) an independent 4-basis for  $S_{\rightarrow}$ , (ii) that  $S_{\rightarrow} = RM_{\rightarrow}$  (thus, a 4-basis for  $RM_{\rightarrow}$ ); and (iii) that  $RM_{\rightarrow}$  can be axiomatized by adding the following “unintelligible” 21-symbol formula to  $R_{\rightarrow}$ :

$$C\ C\ C\ C\ C\ p\ q\ q\ p\ r\ C\ C\ C\ C\ C\ q\ p\ p\ q\ r\ r$$

### New Bases for $RM_{\rightarrow}$ (Cont'd)

- In other words, Meyer & Parks gave the following 5-basis for  $RM_{\rightarrow}$ :

$$C\ p\ p \qquad C\ p\ C\ C\ p\ q\ q \qquad C\ C\ p\ q\ C\ C\ r\ p\ C\ r\ q \qquad C\ C\ p\ C\ p\ q\ C\ p\ q \\ C\ C\ C\ C\ C\ p\ q\ q\ p\ r\ C\ C\ C\ C\ C\ q\ p\ p\ q\ r\ r$$

- The reflexivity axiom  $C\ p\ p$  is dependent in the above 5-basis. The remaining (independent) 4-basis is the Meyer-Parks basis for  $RM_{\rightarrow}$ .
- After much effort (and, with valuable assistance from Bob Veroff and Larry Vos), we (Ernst, Fitelson, Harris) discovered the following 13-symbol replacement for Parks' 21-symbol formula (& there are none shorter [6]):

$$C\ C\ C\ C\ C\ p\ q\ r\ C\ q\ p\ r\ r$$

- The contraction axiom  $C\ C\ p\ C\ p\ q\ C\ p\ q$  is dependent in our new 4-basis. The remaining (independent) 3-basis for  $RM_{\rightarrow}$  contains 31 symbols and 14  $C$ 's (the Meyer-Parks basis has 4 axioms, 48 symbols, and 22  $C$ 's):

$$C\ p\ C\ C\ p\ q\ q \qquad C\ C\ p\ q\ C\ C\ r\ p\ C\ r\ q \qquad C\ C\ C\ C\ C\ p\ q\ r\ C\ q\ p\ r\ r$$

## (Long) Single Axioms for Some Non-Classical Logics

- It was shown by Rezuş [38] (building on earlier seminal work of Tarski and Łukasiewicz [18]) that the systems  $E_{\rightarrow}$ ,  $R_{\rightarrow}$ , and  $L_{\rightarrow}$  have single axioms. However, applying the methods of [38] yields very long, *inorganic* single axioms. As far as we know, these axioms have never been explicitly written down. Here is a 69-symbol (17-variable!) single axiom for the implicational fragment of Łukasiewicz's infinite-valued logic  $L_{\rightarrow}$  (obtained by Ken Harris, using the methods of [38]):

$L_{\rightarrow}$ : *CCCfCgfCCCCCCCdCCeCedCCaCbazzCCCCxyyCCyxwwCCCCtuCutCutssCCqCrqpp*

- Single axioms of comparable length (*i.e.*, containing fewer than 75 symbols) can also be generated for the relevance logics  $E_{\rightarrow}$  and  $R_{\rightarrow}$  (omitted). Here's what we know about the shortest single axioms for the systems  $E_{\rightarrow}$ ,  $R_{\rightarrow}$ ,  $L_{\rightarrow}$ , and  $RM_{\rightarrow}$ :
  - The shortest single axiom for  $E_{\rightarrow}$  has between 23 and 75 symbols.
  - The shortest single axiom for  $R_{\rightarrow}$  has between 23 and 75 symbols.
  - The shortest single axiom for  $L_{\rightarrow}$  has at most 69 symbols.
  - The shortest single axiom for  $RM_{\rightarrow}$  (if there is one<sup>a</sup>) has at least 23 symbols.

<sup>a</sup>Methods of [18] and [38] do *not* apply to  $RM_{\rightarrow}$ , so whether  $RM_{\rightarrow}$  has a single axiom remains open.

## Our "Procedure" for Finding Elegant Axiomatizations

- First, we wrote computer programs to generate a large list of candidate formulas which were to be tested as axioms. For most problems, it was practical to generate an exhaustive list of all formulas with up to twenty-one symbols (we can now do these through 23 symbols).
- All the formulas in the list would be tested (using matrices) to see which are likely to be tautologies in the system in question.<sup>a</sup>
- We immediately eliminated large numbers of formulas by applying known results about axiomatizations in the various systems. For example, as reported by Lemmon et. al., every axiomatization for C5 must contain a formula with  $Cpp$  as a (possibly improper) subformula [16]. Another useful

<sup>a</sup>We say 'likely to be tautologies' because C4 and C5 do not have finite characteristic matrices. Thus, we used matrices which validate all tautologies for the system, but also validate a small number of contingent formulas. This approach works surprisingly well in many systems.

result for this purpose is the Diamond-McKinsey theorem that no Boolean algebra can be axiomatized by formulas containing less than three distinct propositional letters [1, p. 83].

- A set of formulas was selected from the list at random. Using either SEM [56] or a program written by the authors, we found a matrix that respects modus ponens, invalidates a known axiom-basis for the system, but validates the formulas selected from the list. Such a model suffices to show that the formulas are not single axioms for the system.
- All the remaining formulas in the list were then tested against that matrix. Every formula validated by that matrix would be eliminated.
- Steps 4 and 5 were repeated until the list of candidate formulas was down to a small number, or eliminated entirely.
- We would then use OTTER [20] (+strategies!) to attempt to prove a known axiom basis from each of the remaining candidates.

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