

It is customary to setup probability calculi in such a way that:

- Logical/Semantic notions are *presupposed*.
- Conditional Probability is *defined in terms of* unconditional probability (*via* a Ratio Formula).

Suppose we have a function $\text{Pr} : \mathcal{L} \rightarrow \mathbb{R}$ — where \mathcal{L} is a sentential language having connectives \wedge, \vee, \neg — satisfying the following four constraints, for all statements p, q of \mathcal{L} .

- (1) $0 \leq \text{Pr}(p) \leq 1$
- (2) If p and q are *equivalent*, then $\text{Pr}(p) = \text{Pr}(q)$.
- (3) $\text{Pr}(\neg p) = 1 - \text{Pr}(p)$
- (4) $\text{Pr}(p \vee q) = \text{Pr}(p) + \text{Pr}(q)$, if p and q are *incompatible*.

We then *define* conditional probability $\text{Pr}(\cdot | \cdot)$ in terms of unconditional probability $\text{Pr}(\cdot)$, *via* the following Ratio Formula.

$$(5) \quad \text{Pr}(p | q) \stackrel{\text{def}}{=} \frac{\text{Pr}(p \wedge q)}{\text{Pr}(q)} \text{ if } \text{Pr}(q) > 0.$$

Interestingly, (1)–(5) can be adapted to handle both classical, bivalent logic/semantics and trivalent logic/semantics where statements p of \mathcal{L} can be true, false, or neither [2, 4].

The first trick is to add the following clarification, where $\text{Tr}(p)$ asserts that p is true, and $\text{TV}(p)$ asserts that p has a truth-value.

$$(6) \quad \text{Pr}(p) := \frac{\text{Pr}(\text{Tr}(p))}{\text{Pr}(\text{TV}(p))}, \text{ if } \text{Pr}(\text{TV}(p)) > 0.$$

The second trick is to define *equivalence* and *incompatibility* in the following way (in both bivalent and trivalent semantics).

- p and q are *equivalent* if every assignment that makes p true (false) makes q true (false) and *vice versa*.
- p and q are *incompatible* if there is no assignment where both p and q are true.

One reason to favor a trivalent approach is that it can accommodate the addition of an Adams-style conditional — without being subject to Lewisian triviality results [10, 4].

To be more precise, suppose we add a conditional connective (\rightarrow) to our language \mathcal{L} , and we want it to obey the following strong form of Adams's Thesis, for all p, q, x in \mathcal{L} .

$$\mathbf{Adams.} \quad \text{Pr}(p \rightarrow q | x) = \text{Pr}(q | p \wedge x)$$

Adding **Adams** to classical probability calculus yields triviality [7]. *Not so* on the trivalent approach. Indeed, on the trivalent approach, **Adams** is a *theorem*, for bivalent p, q, x [10, 4].

This is because we have the following *modified version* of the Ratio formula for trivalent conditional probability.

$$(7) \quad \text{Pr}(z | x) \stackrel{\text{def}}{=} \frac{\text{Pr}(\text{Tr}(z \wedge x))}{\text{Pr}(\text{TV}(z \wedge x) \wedge \text{Tr}(x))} \text{ if } \text{Pr}(\text{TV}(z \wedge x) \wedge \text{Tr}(x)) > 0.$$

Specifically, if $z = p \rightarrow q$, then — for bivalent p, q, x — we have

Adams (trivalent). For all bivalent p, q, x in \mathcal{L} ,

$$\text{Pr}(p \rightarrow q | x) = \frac{\text{Pr}(\text{Tr}((p \rightarrow q) \wedge x))}{\text{Pr}(\text{TV}((p \rightarrow q) \wedge x) \wedge \text{Tr}(x))} = \frac{\text{Pr}(p \wedge q \wedge x)}{\text{Pr}(p \wedge x)}.$$

Suppose, instead of (1)–(4), we adopt the following four axioms.

- (I) $\text{Pr}(p) = 0$ if $p \vdash \perp$.
- (II) $\text{Pr}(p) = 1$ if $\top \vdash p$.
- (III) If $p \vdash q$, then $\text{Pr}(p) \leq \text{Pr}(q)$.
- (IV) $\text{Pr}(p) + \text{Pr}(q) = \text{Pr}(p \vee q) + \text{Pr}(p \wedge q)$.

Here, we allow “ \vdash ” to be interpreted either in terms of *classical* logical entailment \vdash_C or *intuitionistic* logical entailment \vdash_I .¹

In this way — together with (6) our ratio def. of $\text{Pr}(\cdot | \cdot)$ — we can obtain either classical or intuitionistic probability calculus [21].²

Like classical PC, Intuitionistic PC is in tension with **Adams**. Adding **Adams** to intuitionistic probability calculus yields

No Disconfirmation. For all e, h, k in \mathcal{L} ,

$$\text{Pr}(h | k \wedge e) \geq \text{Pr}(h | k).$$

¹Paraconsistent probability can also be explicated in the (I)–(IV) style [1].

²Some authors in the “Traditional Approach” camp (most notably [22]) take *updating* — rather than *supposing* — to be fundamental for $\text{Pr}(\cdot | \cdot)$.

In fact, we only need the following three general principles about conditional probability (+ associativity and commutativity of \wedge) in order to derive **No Disconfirmation** from **Adams** [16].

$$(C1) \ 0 \leq \Pr(p \mid q) \leq 1$$

$$(C2) \ \Pr(p \mid p \wedge q) = 1$$

$$(C3) \ \Pr(p \wedge q \mid r) = \Pr(p \mid q \wedge r) \cdot \Pr(q \mid r)$$

So, in order to avoid **No Disconfirmation**, we must reject at least one of $\{(C1), (C2), (C3), \text{Adams}\}$. (C1) is simply a choice of finite scale for conditional probability. So, rejecting (C1) seems absurd.

Rejecting (C2) also seems absurd. This would be tantamount to (for some statements p and q) supposing $p \wedge q$ and — under that very supposition — *accepting* “*might not p .*” [24, 25]

It would seem that the only viable option for rejecting **No Disconfirmation** while accepting **Adams** is to reject (C3).

This is what the trivalent approach [10, 4], and two other recent approaches to semantics & probability of indicatives do [3, 8].

Popper [18, 19, 12] proposes an alternative approach, which

- Does *not* presuppose Logical/Semantic notions.
- Does *not* define Conditional Probability in terms of unconditional probability.

In the classical case, Popper [19] gives the following six axioms for conditional probability — as a function $\Pr: \mathcal{L} \times \mathcal{L} \mapsto \mathbb{R}$.

Nontrivial. $(\exists p)(\exists q) [\Pr(p \mid q) \neq \Pr(p \mid p)]$

Identity. $\Pr(p \mid p) \leq \Pr(q \mid q)$

Monotony. $\Pr(p \wedge q \mid r) \leq \Pr(p \mid r)$

Product [(C3)]. $\Pr(p \wedge q \mid r) = \Pr(p \mid q \wedge r) \cdot \Pr(q \mid r)$

Sum. $\Pr(p \mid r) + \Pr(q \mid r) = \Pr(p \wedge q \mid r) + \Pr(p \vee q \mid r)$

Negation. $(\exists p) [\Pr(p \mid r) \neq \Pr(r \mid r)] \implies [\Pr(q \mid r) + \Pr(\neg q \mid r) = \Pr(r \mid r)]$

Popper calls this system \mathbb{B} . All the standard laws of classical conditional probability [*e.g.*, (C1)–(C3)] follow from \mathbb{B} [18, 19] — *without assuming anything about the logic/semantics of \mathcal{L} .*

Popper defines the following relation between statements of \mathcal{L} .

$$p \preceq q \stackrel{\text{def}}{=} (\forall r) [\Pr(p \mid r) \leq \Pr(q \mid r)]$$

$$p \simeq q \stackrel{\text{def}}{=} (p \preceq q) \ \& \ (q \preceq p)$$

\mathbb{B} implies that \simeq is an equivalence relation, which (a) *induces a Boolean algebra \mathcal{B} on \mathcal{L} (viz., the one determined by \models), and (b) is such that $p \simeq q$ ensures *full inter-substitutivity* of p and q .*

In this sense, we can see classical logic/semantics as *emerging* from constraints on conditional probability [6].

If we add **Adams** to \mathbb{B} , then $p \rightarrow q \simeq p \supset q$, and $\Pr(\cdot \mid \cdot)$ is 2-valued [11, 20], *i.e.*, for all p, q in \mathcal{L} , $\Pr(p \mid q) = 0$ or $\Pr(p \mid q) = 1$.

Next, I will examine several weakenings of Popper’s classical theory \mathbb{B} . These can be used to formulate autonomous probability calculi for various subsystems of classical logic.³

³In addition to logics *weaker* than classical logic, the Popperian Approach can also be applied to *extensions/strengthenings* of classical logic [14].

The system \mathbb{M} is given by the following six axioms.

Nonzero. $(\exists p)(\exists q) [\Pr(p \mid q) \neq 0]$

Downbound. $0 \leq \Pr(p \mid q)$

Upbound. $\Pr(p \mid q) \leq \Pr(r \mid r)$

Monotony. $\Pr(p \wedge q \mid r) \leq \Pr(p \mid r)$

Product. $\Pr(p \wedge q \mid r) = \Pr(p \mid q \wedge r) \cdot \Pr(q \mid r)$

\mathbb{M} implies that \simeq is an equivalence relation, which (a) *induces a semilattice structure \mathcal{M} on \mathcal{L} (with \wedge as meet), and (b) is such that $p \simeq q$ ensures *full inter-substitutivity* of p and q .⁴*

$\mathbb{M} + \text{Adams} \implies \text{No Disconfirmation}$ — see Extras slides.

$\mathbb{M} + \text{Sum} = \mathbb{D}$, where \simeq induces a *distributive lattice \mathcal{D} on \mathcal{L}* .⁵

⁴In subsystems of \mathbb{B} , we also need $p \simeq q \implies (\forall r) [\Pr(r \mid p) = \Pr(r \mid q)]$, so as to ensure $p \simeq q$ guarantees inter-substitutivity of p and q on the right-hand sides of $\Pr(\cdot \mid \cdot)$ ’s — see page 10 of [19] for discussion.

⁵Narens [17] takes a Traditional Approach to lattice-theoretic probability.

If we add the following two axioms to \mathbb{D} , then \simeq induces a *distributive lattice with zero/bottom element* (\perp) \mathcal{D}_\perp on \mathcal{L} . I will call the resulting probability calculus \mathbb{D}_\perp .

Zero₁. $\Pr(p | q) \leq \Pr(q | \perp)$

Zero₂. $[\Pr(p | q) \neq \Pr(p | p)] \implies [\Pr(p | \perp) + \Pr(\perp | q) = \Pr(p | p)]$

In order to get from \mathbb{D}_\perp to intuitionistic probability calculus \mathbb{H} (with intuitionistic conditional \multimap), we need two more axioms.

Conditional₁. $\Pr(p \multimap q | p \wedge r) = \Pr(q | p \wedge r)$

Conditional₂. $(p \wedge q) \leq r \implies p \leq (q \multimap r)$

Then, the algebra induced by \simeq on \mathcal{L} is a Heyting Algebra \mathcal{H} [12]. And, if we add **Negation** to \mathbb{H} , this brings us back to \mathbb{B} .

The axiomatization of intuitionistic probability calculus remains controversial. Some authors assume **Adams** rather than **Conditional₁** [15], which leads to **No Disconfirmation** [16, 21].

In [16, 15], $\Pr(h | k) \leq \Pr(h | e \wedge k)$ is interpreted as implying only “if we can *prove* h from k , we can *prove* h from $e \wedge k$.”

On this interpretation, **No Disconfirmation** seems benign. But, as Weatherson [21] points out, a *Bayesian* with “intuitionsitic tendencies” may reasonably reject **No Disconfirmation**.

They may interpret $\Pr(h | k) \leq \Pr(h | e \wedge k)$ as implying that *no e can disconfirm any h , relative to any background corpus k* .

On either the Traditional or the Popperian approach, we need to be clear on how we are understanding $\Pr(p | q)$. I follow Weatherson in taking a Bayesian stance here.

That is, I prefer to understand $\Pr(p | q)$ as measuring a (*rational*) *degree of belief in p , on the (indicative) supposition that q* [5].

The “semantic/logical” interpretation of $\Pr(\cdot | \cdot)$ is less clear to me. Some authors seem to interpret $\Pr(p | q)$ as “the degree to which q is deducible from p ” [11]. And, some have an externalist epistemic (“evidential probability”) interpretation [23].

I will close with three questions that I’m hoping to answer.

- Can the Popperian Approach be used for the trivalent case? That is, can we write down axioms for $\Pr(\cdot | \cdot)$ which
 - (a) give us the properties of $\Pr(\cdot | \cdot)$ we seek (*e.g.*, **Adams** without triviality or **No Disconfirmation**), and
 - (b) induce (*via* \simeq) the semantical structures of, *e.g.*, [2, 4].
- Can the Popperian Approach be applied to other non-classical systems that have been developed recently — especially in connection with indicative conditionals and epistemic modals. For instance, the systems being explored by Goldstein & Santorio [8] and Holliday & Mandelkern [9]?
- More generally, I like the idea of taking *conditional probability judgments regarding pairs of sentences* (of \mathcal{L}) as fundamental [3]. It would be nice if we could write down axioms for *that*, and let *those* determine the underlying logic/semantics for \mathcal{L} — *à la* Popper. But, *can this be done*?

Here is a proof (sketch) of **No Disconfirmation** in Popper’s system \mathbb{M} . We will first prove (C1), (C2), and a **Lemma**.

(C1) $0 \leq \Pr(p | q) \leq 1$

Proof.

We just need to prove $\Pr(p | q) \leq 1$. In light of **Upbound**, it suffices to show that $\Pr(p | p) = 1$. By **Upbound**, $\Pr(p | p)$ is independent of p . So, set $\Pr(p | p) = k$ for all p . Then, from **Downbound** and **Monotony**, we have

$$\Pr(p | p \wedge p) \leq k = \Pr(p \wedge p | p \wedge p) \leq \Pr(p | p \wedge p)$$

$$\Pr(p | p \wedge (p \wedge p)) \leq k = \Pr(p \wedge (p \wedge p) | p \wedge (p \wedge p)) \leq \Pr(p | p \wedge (p \wedge p))$$

$$\therefore \Pr(p | p \wedge p) = \Pr(p | p \wedge (p \wedge p)) = k$$

Then, by **Product**, we have

$$k = \Pr(p \wedge p | p \wedge p) = \Pr(p | p \wedge (p \wedge p)) \cdot \Pr(p | p \wedge p) = k^2$$

Hence, either $k = 0$ or $k = 1$. By **Product**, **Upbound**, and **Monotony** ($\times 2$)

$$k \cdot \Pr(x | z) = \Pr(x \wedge z | x \wedge z) \cdot \Pr(x | z)$$

$$= \Pr((x \wedge z) \wedge x | z) \leq \Pr(x \wedge z | z) \leq \Pr(x | z) \leq k$$

Thus, if $k = 0$, then $\Pr(x | z) = 0$ for all x, z , which contradicts **Nonzero**. \square

$$(C2) \Pr(x | x \wedge z) = 1 = \Pr(x | z \wedge x)$$

Proof.

Since by (C1) the value of $\Pr(x | z \wedge y)$ lies in $[0, 1]$, we obtain the following inequality immediately from **Product**

$$(i) \quad \Pr(x \wedge z | y) \leq \Pr(z | y)$$

By $\Pr(p | p) = 1$, **Monotony**, and **Upbound**, we have

$$(ii) \quad 1 = \Pr(x \wedge z | x \wedge z) \leq \Pr(x | x \wedge z) \leq 1$$

and likewise, using (i) in place of **Monotony**, we have

$$(iii) \quad 1 = \Pr(z \wedge x | z \wedge x) \leq \Pr(x | z \wedge x) \leq 1$$

Together, (ii) and (iii) imply (C2). \square

Lemma. If $\Pr(x | y \wedge z) = 1$, then $\Pr(y | z) \leq \Pr(x | z)$.

Proof.

By **Product**, we have

$$(iv) \quad \Pr(x | y \wedge z) \cdot \Pr(y | z) = \Pr(x \wedge y | z)$$

Suppose $\Pr(x | y \wedge z) = 1$.

From (iv) and **Monotony**, we may deduce

$$\Pr(y | z) = \Pr(x \wedge y | z) \leq \Pr(x | z)$$

Therefore, $\Pr(y | z) \leq \Pr(x | z)$. \square

No Disconfirmation. $\Pr(h | k \wedge e) \geq \Pr(h | k)$

Proof.

$$(v) \Pr(h | e \wedge (h \wedge k)) = 1$$

- by (C2) + associativity & commutativity of \wedge .

$$(vi) \Pr(e \rightarrow h | h \wedge k) = 1$$

- by (v) + **Adams**.

$$(vii) \Pr(h | k) \leq \Pr(e \rightarrow h | k)$$

- by (vi) + **Lemma** (where $x := e \rightarrow h$, $y := h$, $z := k$).

$$(viii) \Pr(h | k) \leq \Pr(h | k \wedge e)$$

- by (vii) + **Adams** + commutativity of \wedge .

Popper and Miller [19] give a formal proof of the associativity and commutativity of \wedge from the axioms of \mathbb{M} (+ the assumption that $p \simeq q \Rightarrow (\forall r) [\Pr(r | p) = \Pr(r | q)]$ — see footnote 4). \square

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