

Another Argument for \supset

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In this note, I will outline an argument for \supset , which combines (and refines) arguments from Priest and Gibbard. First, some notation. I will use ' \rightarrow ' for the English indicative conditional, and ' \supset ' for material implication. I will (typically) use English connectives in the meta-theory (sometimes, I will abbreviate the meta-theoretic conditional as ' \Rightarrow ', and sometimes I will use the word "entails", which is meant to be synonymous with the meta-theoretic \Rightarrow), and I will assume that our meta-theory is *classical*. I will give an argument for the following meta-theoretic statement (understood as a *schema*, which holds for *any* p/q):

$$p \rightarrow q \text{ if and only if } p \supset q \\ \text{[i.e., } p \rightarrow q \Leftrightarrow p \supset q \text{]}$$

This requires establishing the following two meta-theoretic conditionals:

- ① If $p \rightarrow q$, then $p \supset q$. [i.e., $p \rightarrow q \Rightarrow p \supset q$]
- ② If $p \supset q$, then $p \rightarrow q$. [i.e., $p \supset q \Leftarrow p \rightarrow q$]

My strategy will be to prove ① first, and then *use* ① to prove ②. Here goes.

Argument for ①. Assuming a classical meta-theory, ① requires *only* the following principle:

(MP₋) If p and $p \rightarrow q$, then q . [i.e., *Modus Ponens* for ' \rightarrow ' *preserves truth*.]

Here is my argument for ①. I will actually prove the *contrapositive* of ①.

1	$p \supset q$ is false.	Assumption (for \Rightarrow I)						
2	p is true.	From (1), by classical logic.						
3	q is false.	From (1), by classical logic.						
4	<table style="border-collapse: collapse; width: 100%;"> <tr> <td style="border-left: 1px solid black; padding-left: 10px;">$p \rightarrow q$ is true.</td> <td style="padding-left: 20px;">Assumption (for RAA)</td> </tr> <tr> <td style="border-left: 1px solid black; padding-left: 10px;">q is true.</td> <td style="padding-left: 20px;">From (2) and (4), by (MP₋).</td> </tr> <tr> <td style="border-left: 1px solid black; padding-left: 10px;">Contradiction.</td> <td style="padding-left: 20px;">From (3), (5).</td> </tr> </table>	$p \rightarrow q$ is true.	Assumption (for RAA)	q is true.	From (2) and (4), by (MP ₋).	Contradiction.	From (3), (5).	
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q is true.	From (2) and (4), by (MP ₋).							
Contradiction.	From (3), (5).							
7	$p \rightarrow q$ is false.	From (4)–(6), by (RAA).						
8	$p \supset q$ is false \Rightarrow $p \rightarrow q$ is false.	From (1)–(7), by (\Rightarrow I).						
9	①	From (8), by \Rightarrow contraposition. \square						

Thus, assuming a classical meta-theory, *all* we need in order to prove ① is (MP₋). This shows that ① is *virtually equivalent to the assertion that Modus Ponens is truth-preserving for the indicative conditional*.

Argument for ②. My argument for ② depends on the following six principles:

(EXP₋) If ' $(p \& q) \rightarrow r$ ', then ' $p \rightarrow (q \rightarrow r)$ '. [i.e., *Exportation* for ' \rightarrow ' *preserves truth*.]

① If $p \rightarrow q$, then $p \supset q$.

(AND₋) ' $(p \& q) \rightarrow q$ ' is a logical truth.

(LTE) If p is a logical truth, and p entails q , then q is a logical truth.

(SUB) If p' is obtained from p by substitution of logical equivalents (i.e., if p' results from substituting q' for q in p , where $q' \Leftrightarrow q$), then p entails p' .

(SDT₃¹) If ' $p \supset q$ ' is a logical truth, then p entails q .

Here is my argument for ②. This argument will be *direct*.

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| 1 | ‘ $(p \& q) \rightarrow q$ ’ is a logical truth. | (AND ₋) |
| 2 | ‘ $((p \supset q) \& p) \rightarrow q$ ’ is a logical truth. | From (1), by (SDT ₃ ¹), (SUB), and (LTE). |
| 3 | ‘ $(p \supset q) \rightarrow (p \rightarrow q)$ ’ is a logical truth. | From (2), by (EXP ₋) and (LTE). |
| 4 | ‘ $(p \supset q) \supset (p \rightarrow q)$ ’ is a logical truth. | From (3), by ① and (LTE). |
| 5 | ② | From (4), by (SDT ₃ ¹). |

Therefore, *the only way* one can resist the conclusion that the English indicative conditional \rightarrow is *equivalent* to \supset is to reject some of the following six assumptions (or some other classical inference in the meta-theory):

(MP₋) If p and $p \rightarrow q$, then q . [*i.e.*, *Modus Ponens* for ‘ \rightarrow ’ *preserves truth*.]

(EXP₋) If ‘ $(p \& q) \rightarrow r$ ’, then ‘ $p \rightarrow (q \rightarrow r)$ ’. [*i.e.*, *Exportation* for ‘ \rightarrow ’ *preserves truth*.]

(AND₋) ‘ $(p \& q) \rightarrow q$ ’ is a logical truth.

(LTE) If p is a logical truth, and p entails q , then q is a logical truth.

(SUB) If p' is obtained from p by substitution of logical equivalents (*i.e.*, if p' results from substituting q' for q in p , where $q' \Leftrightarrow q$), then p entails p' .

(SDT₃¹) If ‘ $p \supset q$ ’ is a logical truth, then p entails q .

As we have seen, MacFarlane & Kolodny and McGee reject (MP₋). McGee seems to accept *all the other* assumptions of this argument, whereas MacFarlane & Kolodny also reject (EXP₋). I think MacFarlane & Kolodny accept everything here *except* for (MP₋) and (EXP₋). But, it’s not at all obvious to me why someone who’s worried about \supset should accept (SDT₃¹). That places a non-trivial constraint on our meta-theoretic “entailment” and “equivalence” relations, which could “trickle down” to our indicative, *especially* if we were inclined to assume some sort of semantic deduction theorem(s) for our *indicative* conditional as well. Consider:

(SDT₁¹) If ‘ $p \rightarrow q$ ’ is a logical truth, then p entails q .

(SDT₂²) If p entails q , then ‘ $p \rightarrow q$ ’ is a logical truth.

(SDT₃²) If p entails q , then ‘ $p \supset q$ ’ is a logical truth.

I suspect that (SDT)-type assumptions are doing more work than meets the eye here, since it is easy to tacitly presuppose that (SDT)-type principles hold for *both* \rightarrow *and* \supset . But, of course, if (SDT) *does* hold for *both* connectives, then we can “prove” a *validity-preserving* rendition of the desired equivalence, *trivially*:

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|---|---|--|
| 1 | ‘ $p \supset q$ ’ is a logical truth. | Assumption (for \Rightarrow I) |
| 2 | p entails q . | From (1), by (SDT ₃ ¹). |
| 3 | ‘ $p \rightarrow q$ ’ is a logical truth. | From (2), by (SDT ₂ ²). |
| 4 | $\models p \supset q \Rightarrow \models p \rightarrow q$ | From (1)–(3), by \Rightarrow I. \square |

and

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|---|---|--|
| 1 | ‘ $p \rightarrow q$ ’ is a logical truth. | Assumption (for \Rightarrow I) |
| 2 | p entails q . | From (1), by (SDT ₁ ¹). |
| 3 | ‘ $p \supset q$ ’ is a logical truth. | From (2), by (SDT ₃ ²). |
| 4 | $\models p \rightarrow q \Rightarrow \models p \supset q$ | From (1)–(3), by \Rightarrow I. \square |

And, if *these* meta-theoretic claims can be shown, then my conjecture about \rightarrow and \supset having similar *validity-preserving* (pure conditional) forms starts to look more plausible. We need to think more about (SDT)’s.