## Sample In-Class Final

Philosophy 12A May 6, 2010

This is a sample in-class final examination. The structure of the actual in-class final will be exactly the same as this sample. The particular problems on the actual final will be different. On the actual final exam, you can expect to see problems of roughly the same difficulty as these (perhaps, slightly harder).

## 1 Ten Cumulative True/False Questions

- **T F** 1. If it is not possible for the conclusion of an argument to be false, then the argument is valid.
- **T F** 2. All sound arguments are valid.
- **T F** 3. Some valid arguments are unsound.
- T F 4. If the conclusion of a valid argument is true, the premises must be true as well.
- **T F** 5. The following argument is sound (absolutely): "If Prince William is unmarried, then Prince William is a bachelor. Prince William is a bachelor. Therefore, Prince William is unmarried."
- **T F** 7. The following LMPL argument is valid:  $(\forall x)(\exists y)(Fx \to Gy) \\ \therefore (\exists x)(\forall y)(Fx \to Gy)$
- **T F** 8. The following is a legitimate application of the  $\forall$ I rule:
  - 1
     (1)  $(\forall x)(\mathsf{Fx} \rightarrow \mathsf{Gx})$  Premise

     2
     (2) Fb
     Premise

     1
     (3)  $\mathsf{Fb} \rightarrow \mathsf{Gb}$  1
      $\forall \mathsf{E}$  

     1,2
     (4) Gb
     2,3
      $\rightarrow \mathsf{E}$  

     1,2
     (5)  $(\forall x)\mathsf{Gx}$  4
      $\forall \mathsf{I}$
- **T F** 9. The following is a legitimate application of the  $\exists E$  rule:
  - 1 (3x)Fx Premise (1) 2 Premise (2) Ga 3 (3) Fa Assumption 2,3 Fa&Ga 2,3 &I (4) 4 3I (5)  $(\exists x)(Fx\&Gx)$ (6)  $(\exists x)(Fx\&Gx)$ 1,3,5 3E
- **T F** 10. The following LSL sentence is a *tautology*:  $(A \lor B) \leftrightarrow ((A \to B) \to B)'$ .

**Answers**: (1) **T**, (2) **T**, (3) **T**, (4) **F**, (5) **F**, (6) **T**, (7) **F**, (8) **F**, (9) **F**, (10) **T**.

#### 2 An LSL Semantics Problem

Do a full truth-table for the following LSL sentence:  $(A \lor B) \leftrightarrow ((A \to B) \to B)'$ .

[I'll leave the completion of this full truth-table as an exercise.]

#### 3 Another LSL Semantics Problem

Determine whether the following statement is correct, using any legitimate truth-table technique.

$$A \& (B \rightarrow C) \vDash (A \& C) \lor (A \& \sim B)$$

[I'll leave the solution to this LSL semantics problem as an exercise.]

#### 4 Completing an Incomplete LMPL Natural Deduction

Complete the following natural deduction proof of the LMPL theorem  $\vdash (\forall x)[(Fx \rightarrow Gx) \lor (Gx \rightarrow Fx)]$ . All assumptions and premises have been filled in. Fill in any missing labels (right) and premise/assumption numbers (left). NOTE: *SI* and *TI* might be used in this proof.

		(1)	$Ga \lor \sim Ga$	
	2	(2)	Ga	Assumption
		(3)	$Fa \rightarrow Ga$	
		(4)	$(Fa \rightarrow Ga) \lor (Ga \rightarrow Fa)$	
_	5	(5)	~Ga	Assumption
		(6)	$Ga \rightarrow Fa$	
		(7)	$(Fa \rightarrow Ga) \lor (Ga \rightarrow Fa)$	
_		(8)	$(Fa \rightarrow Ga) \lor (Ga \rightarrow Fa)$	
_		(9)	$(\forall x)[(Fx \to Gx) \lor (Gx \to Fx)]$	<b>*</b>

Answer:

	(1)	$Ga \lor \sim Ga$	TI (LEM)
2	(2)	Ga	Assumption
2	(3)	$Fa \rightarrow Ga$	2 SI (PMI)
2	(4)	$(Fa \rightarrow Ga) \lor (Ga \rightarrow Fa)$	3 vI
5	(5)	~Ga	Assumption
5	(6) $Ga \rightarrow Fa$	Ga → Fa	5 SI (PMI)
5	(7)	$(Fa \rightarrow Ga) \lor (Ga \rightarrow Fa)$	6 ∨I
	(8)	$(Fa \rightarrow Ga) \lor (Ga \rightarrow Fa)$	1, 2, 4, 5, $7 \lor E$
	(9)	$(\forall x)[(Fx \to Gx) \lor (Gx \to Fx)]$	8 ∀I <b>◆</b>

# 5 Diagnosing an Incorrect LMPL Natural Deduction

Explain what is wrong with the following *alleged* proof of the sequent  $(\forall x)(\exists y)(Fx\&Gy) \vdash (\exists y)(\forall x)(Fx\&Gy))$ . Is it possible to *fix* this *illegitimate* proof? If so, explain how. If not, explain why it *can't* be fixed. Try to find a counterexample LMPL interpretation. If you can't find one that contains *fewer than 3 individuals* in it, then you can assume that the argument is *valid*, and that it *can* be proven.

1	(1)	$(\forall x)(\exists y)(Fx \& Gy)$	Premise
1	(2)	$(\exists y)(Fa \& Gy)$	1 ∀E
3	(3)	Fa & Gb	Assumption
3	(4)	$(\forall x)(Fx \& Gb)$	3 ∀I
3	(5)	$(\exists y)(\forall x)(Fx \& Gy)$	4 ∃I
1	(6)	$(\exists v)(\forall x)(Fx \& Gv)$	2. 3. 5 ∃E

**Answer**: The problem with this proof occurs at step (4). The constant 'a' occurs in line (3) upon which line (4) depends. This violates the restriction on t in the definition of the  $\forall$ I rule. The sequent is valid, however, (if you try to find a counterexample interpretation, you'll quickly be forced to have *more than 2* individuals in your domain). I proved this LMPL sequent in 12 steps in my lecture notes.

#### 6 Working with a Given LMPL Interpretation

Consider the following LMPL interpretation:

Explain why  $\mathcal{I}$  shows that the following claim is true:

$$(\forall x)(Fx \to Gx) \to (\forall x)(Hx \to Jx) \nvDash (\exists x)(Fx \& Gx) \to (\forall x)(Hx \to Jx)$$

**Answer**: The premise ' $(\forall x)(Fx \to Gx) \to (\forall x)(Hx \to Jx)$ ' is true, because its antecedent ' $(\forall x)(Fx \to Gx)$ ' is false, since its instance ' $Fb \to Gb$ ' is false on  $\mathcal{I}[\beta \in \operatorname{Ext}(F), \beta \notin \operatorname{Ext}(G)]$ . The conclusion ' $(\exists x)(Fx \& Gx) \to (\forall x)(Hx \to Jx)$ ' is false. Its antecedent ' $(\exists x)(Fx \& Gx)$ ' is true, since its instance 'Fa & Ga' is true on  $\mathcal{I}[\alpha \in \operatorname{Ext}(F), \alpha \in \operatorname{Ext}(G)]$ . But, its consequent ' $(\forall x)(Hx \to Jx)$ ' is false, because its instance ' $Ha \to Ja$ ' is false on  $\mathcal{I}[\alpha \in \operatorname{Ext}(H), \alpha \notin \operatorname{Ext}(J)]$ . So,  $\mathcal{I}$  shows that this is an invalid argument.

## 7 Constructing a Counterexample LMPL Interpretation

Construct an LMPL interpretation which shows that the following claim is true (make sure to *explain why* your interpretation is a counterexample to the validity of the argument):

$$(\forall x)(\exists y)(Gy \rightarrow Fx) \not\models (\forall x)[(\exists y)Gy \rightarrow Fx]$$

#### 8 An LSL Natural Deduction

Give a natural deduction proof of the following LSL sequent. You may use SI and TI (you don't have to).

$$S \rightarrow (R \vee P), P \rightarrow (\sim R \rightarrow Q) \vdash S \rightarrow (Q \vee R)$$

**Answer**: Here's a 13-step proof, using SI:

1	(1) S→(R∨P) (2) P→(~R→C	Premise Premise
3	(3) S	Assumption
	(4) R <sub>&gt;</sub> P	1,3 →Ė
1,3 5	(5) R	Assumption
5	(6) Q√R	5 VI
7	(7) P	Assumption
2,7	(8) ~R→Q	2,7 →E
2,7	(9) ~~R√Q	8 SI (IMP)
2,7	(10) R <sub>V</sub> Q	9 SI (SDN)
2,7	(11) Q <sub>V</sub> R	10 SI (Com)
1,2,3	(12) Q <sub>V</sub> R	4,5,6,7,11 VE
1,2	$(13) S \rightarrow (Q \lor R)$	3,12 →

#### 9 An LMPL Natural Deduction

Give a natural deduction proof of the following LMPL argument. You may use SI and TI (don't have to).

$$(\forall x)[Lx \to ((\forall y)(Py \to Vy) \to Mx)]$$

$$(\exists z)(Pz \& Vz) \to (\forall y)(Py \to Vy)$$

$$\therefore (\exists x)Lx \to [(\exists z)(Pz \& Vz) \to (\exists y)My]$$

**Answer**: Here's a 13-step proof, without using SI/TI:

1	(1)	$(\forall x)(Lx\rightarrow((\forall y)(Py\rightarrow Vy)\rightarrow Mx))$	Premise
2	(2)	$(\exists z)(Pz\&Vz)\rightarrow (\forall y)(Py\rightarrow Vy)$	Premise
3	(3)	(Jx)Lx	Assumption
4	(4)	(∃z)(Pz&Vz)	Assumption
5	(5)	La	Assumption
1	(6)	La→((∀y)(Py→Vy)→Ma)	1 ∀E
1,5	(7)	(∀y)(Py→Vy)→Ma	6,5 →E
2,4	(8)	(∀y)(Py→Vy)	2,4 →E
1,2,4,5	(9)	Ма	7,8 →E
1,2,4,5	(10)	(∃y)My	IE 9
1,2,3,4	(11)	(Jy)My	3,5,10 <b>3</b> E
1,2,3	(12)	(∃z)(Pz&Vz)→(∃y)My	4,11 →I
1,2	(13)	$(\exists x)Lx \rightarrow ((\exists z)(Pz\&Vz) \rightarrow (\exists y)My)$	3,12 →

## 10 Another LMPL Natural Deduction

Give a natural deduction proof of the following LMPL sequent. You may use SI and TI (don't have to).

$$\sim (\exists x)(\forall y)(Fx \to Fy) \vdash \bot$$

**Answer**: Here's a 12-step proof, using SI:

1	(1)	$\sim$ ( $\exists x$ )( $\forall y$ )( $Fx \rightarrow Fy$ )	Premise
1	(2)	$(\forall x) \sim (\forall y) (Fx \rightarrow Fy)$	SI (QS)
1	(3)	~(∀y)(Fa→Fy)	2 VE
1	(4)	(∃y)~(Fa→Fy)	3 SI (QS)
5	(5)	~(Fa→Fb)	Ass
5	(6)	Fa&~Fb	5 SI (Neg-Imp)
5	(7)	~Fb	6 &E
5	(8)	Fb→Fc	7 SI (PMI)
5	(9)	(∀y)(Fb→Fy)	8 AI
5	(10)	$(\exists x)(\forall y)(Fx \rightarrow Fy)$	IE 0
1,5	(11)	Λ	1,10 ~E
1	(12)	Λ	4,5,11 ∃E