

Homework #2 Solutions

February 22, 2010

Page 19: Exercises 1, 3, 5 & 12.

- E : Van Gogh's pictures are the world's most valuable.
 P : Van Gogh's pictures are the world's most profound.
Final symbolization: $E \& \sim P$.
- Same lexicon as in (1).
Final symbolization: $\sim E \& \sim P$.
- D : Digital computers can simulate every aspect of human intelligence.
 N : Neural networks can simulate every aspect of human intelligence.
 E : Digital computers can simulate some aspects of human intelligence.
 O : Neural networks can simulate some aspects of human intelligence.
Final symbolization: $(\sim D \& \sim N) \& (E \& O)$.
- R : It rains.
 P : It pours.
Final symbolization: $R \rightarrow P$. [also acceptable: anything equivalent, since the form is not obvious from the English sentence — *e.g.*, $\sim(R \& \sim P)$]

Page 26: Exercises 2, 4, 6 & 16.

- T : It is Tuesday.
 B : It is Belgium.
 L : I'm lost.
Final symbolization: $(T \& \sim B) \rightarrow L$.
- D : The economy declines.
 C : There is a change of leadership.
 R : There will be a recession.
Final symbolization: $D \rightarrow (\sim C \rightarrow R)$. [also acceptable: $D \rightarrow (C \vee R)$]
- A : Applicants may examine their dossiers.
 W : Applicants have already waived their right to examine their dossiers.
 R : Applicants' referees approve.
Final symbolization: $A \rightarrow (\sim W \& R)$.
- R : There is a right to smoke in public.
 H : Smoking in public significantly affects the health of others.
Final symbolization: $(R \rightarrow \sim H) \& \sim \sim H$. [note: we want to stay as close to the English *form* as possible.]

Page 33: Exercises 1, 5, 10 & 12.

- G : The government rigs the election.
 R : There will be riots.
 V : The government is guaranteed victory.
Final symbolization (argument): $G \rightarrow R, \sim G \rightarrow V, \sim G \rightarrow \sim V \therefore R$
- K : I know I exist.
 E : I exist.
 H : I think.
 N : I know I think.
Final symbolization (argument): $K \rightarrow E, (N \rightarrow K) \& (H \rightarrow N), H \therefore E$
- N : At least two contestants enter.
 C : There will be a contest.
 W : There will be a winner.
 E : All contestants perform equally well.
 L : There is a loser.
Final symbolization (argument): $\sim N \rightarrow \sim C, \sim C \rightarrow \sim W, E \rightarrow \sim W, \sim L \leftrightarrow \sim W \therefore L \rightarrow (N \& \sim E)$

12. *D*: The Mayor is defeated.
S: Council members are involved in a financial scandal.
U: The urban middle class supports the Mayor.
 Final symbolization (argument): $\sim D \rightarrow U, U \rightarrow \sim S \therefore D \leftrightarrow S$

Page 43, IV: 2 & 6

2. This is an erroneous usage of scare quotes. 'Rome' is not the largest city in Italy, since 'Rome' a *word*, *not a city*.
6. This is a correct usage of selective (corner) quotes. It's a true (metatheoretic) statement. [If we had used scare quotes here instead, then *that* would have been a mistake, since ' $\sim p$ ' is *not* a wff of LSL, because there are only upper case letters in LSL.]

Page 57, I: 1 & 5

1. Truth-Table for ' $A \rightarrow (B \rightarrow (A \& B))$ ' (main connective in red):

A	B	$A \rightarrow (B \rightarrow (A \& B))$
T	T	T
T	\perp	T
\perp	T	\perp
\perp	\perp	T

\therefore ' $A \rightarrow (B \rightarrow (A \& B))$ ' is *tautological* (it is true on all interpretations).

5. Truth-Table for ' $((F \& G) \rightarrow H) \rightarrow ((F \vee G) \rightarrow H)$ ' (main connective in red):

F	G	H	$((F \& G) \rightarrow H) \rightarrow ((F \vee G) \rightarrow H)$
T	T	T	T
T	T	\perp	T
T	\perp	T	\perp
T	\perp	\perp	T
\perp	T	T	T
\perp	T	\perp	T
\perp	\perp	T	T
\perp	\perp	\perp	T

\therefore ' $((F \& G) \rightarrow H) \rightarrow ((F \vee G) \rightarrow H)$ ' is *contingent* (it is true on some interpretations, false on others).

Page 58, II

- II. Truth-Tables for the sentences in question (main connectives in red):

(1)	<table border="1"><thead><tr><th>A</th><th>B</th><th>$A \vee B$</th></tr></thead><tbody><tr><td>T</td><td>T</td><td>T</td></tr><tr><td>T</td><td>\perp</td><td>T</td></tr><tr><td>\perp</td><td>T</td><td>T</td></tr><tr><td>\perp</td><td>\perp</td><td>\perp</td></tr></tbody></table>	A	B	$A \vee B$	T	T	T	T	\perp	T	\perp	T	T	\perp	\perp	\perp	(2)	<table border="1"><thead><tr><th>A</th><th>B</th><th>$A \rightarrow B$</th></tr></thead><tbody><tr><td>T</td><td>T</td><td>T</td></tr><tr><td>T</td><td>\perp</td><td>\perp</td></tr><tr><td>\perp</td><td>T</td><td>T</td></tr><tr><td>\perp</td><td>\perp</td><td>T</td></tr></tbody></table>	A	B	$A \rightarrow B$	T	T	T	T	\perp	\perp	\perp	T	T	\perp	\perp	T	(3)	<table border="1"><thead><tr><th>A</th><th>B</th><th>$\sim (A \& \sim B)$</th></tr></thead><tbody><tr><td>T</td><td>T</td><td>\perp</td></tr><tr><td>T</td><td>\perp</td><td>T</td></tr><tr><td>\perp</td><td>T</td><td>\perp</td></tr><tr><td>\perp</td><td>\perp</td><td>T</td></tr></tbody></table>	A	B	$\sim (A \& \sim B)$	T	T	\perp	T	\perp	T	\perp	T	\perp	\perp	\perp	T
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Therefore, we have the following equivalences:

- (6) and (7) are equivalent.
- (1) and (4) are equivalent.
- (2), (3), and (5) are equivalent.

Page 58, III

- III. No, if p is not a tautology, it does *not* follow that ' $\sim p$ ' is a tautology. This is equivalent to the metatheoretic question: "If $\#p$, then does it follow that $\# \sim p$?" There are LSL sentences such that *both* $\#p$ and $\# \sim p$. Any atomic wff (e.g., ' A ') will do. More generally, any *contingent* sentence p will, by definition, be such that *both* $\#p$ and $\# \sim p$.